## Project Problems

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1. Given an arbitrary graph  $\mathbf{G} = \langle \mathbf{V}, \mathbf{E} \rangle$  with positive and negative edge-weights, what is the minimum number of edges to delete so that the resultant graph does not contain a negative weight cycle?

References and Motivation - It is well known that graphs can be used to represent systems of difference constraints [CLR92]. We are thus trying to minimize the number of constraints to be eliminated to make the system feasible. It is known that the problem is NP-complete.

Project credit - Find a constant bound approximation ( or better! )

2. **Definition:** 0.1 A graph is said to be an interval graph if its vertices can be put into 1-1 correspondence with a set of intervals on the real-line, so that two intervals intersect if and only if their corresponding vertices in the graph have an edge between them.

Given an arbitrary graph  $G = \langle V, E \rangle$  (unweighted) what is the minimum number of edges to delete so that the resultant graph is an interval graph?

**References and Motivation** - Interval graphs are used to model a number of problems in diverse domains [Spi]. The interval property is crucial to the solution of these problems.

**Project Credit** - Show that the problem is NP-complete and find a constant factor approximation algorithm.

3. **Definition:** 0.2 A graph G is said to be a split graph, if its vertex set can be partitioned into two parts  $V_1$  and  $V_2$  such that the subgraph induced by  $V_1$  is an interval graph and the that induced by  $V_2$  is an independent set.

Given an arbitrary graph  $G = \langle V, E \rangle$  (unweighted) is it a split graph?

**References and Motivation** - Split interval graphs are used in modeling problems in Computational Biology [MWZ98]. The hardness of this problem is not known.

**Project credit** - Show that the problem is NP-complete.

- 4. **Definition:** 0.3 A graph G is said to be a probe-interval graph if there exists a probe/non-probe assignment to its vertices i.e. a partition of the vertex set into P, N, such that
  - (a) The graph G is an interval graph, and
  - (b) For every edge, one of the vertices is is in the set **P**?

References and Motivation - Probe interval graphs are used to model gene sequencing problems in Computational Biology [JS00]. It is known that the problem can be solved in polynomial time if the probe/non-probe assignment is given.

**Project Credit** - Show that the problem is NP-complete.

5. **Definition:** 0.4 A graph  $G = \langle V, E \rangle$  is said to be collapsible, if for every subset S of V, with |S| even (i.e. all even subsets of the vertex set), G has a spanning connected subgraph H such that the set of odd degree vertices in H is precisely S.

Given an arbitrary graph **G**, is it collapsible?

References and Motivation - This is a pure Discrete Structure problem, with absolutely no practical relevance! However, establishing the hardness of the problem will settle a 10 year old conjecture [CL93]

Project Credit - Show that the problem is coNP-complete.

## References

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- [CLR92] T. H. Cormen, C. E. Leiserson, and R. L. Rivest. *Introduction to Algorithms*. MIT Press and McGraw-Hill Book Company, 6th edition, 1992.
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- [Spi] J. Spinrad. Graph Theory. http://www.vuse.vanderbilt.edu/spin/research.html.