

A Hybrid-ARQ Protocol Using Noncoherent Orthogonal Modulation

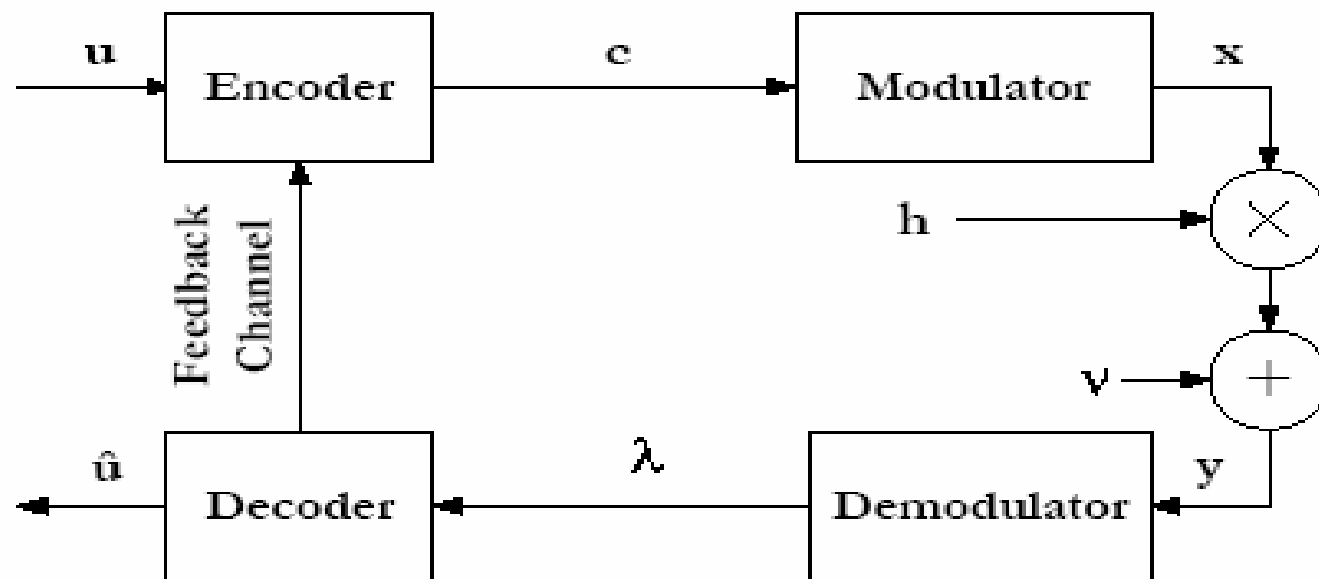
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Overview

- **Noncoherent orthogonal modulation**
 - Energy detection based scheme
 - Pseudo-orthogonal modulation used in Bluetooth, Zigbee etc.
- **Hard Decision Decoding with Reed Solomon codes**
 - Channel Coding and modulation can be made independent
 - Appropriate bound is the Discrete Memoryless Capacity (DMC)
- **Hybrid-ARQ technique**
 - Combines FEC with ARQ
 - Codeword broken into B blocks; incremental redundancy and code-combining
- **Relaying**
 - Exploits spatial diversity from redundant nodes
 - Hybrid-ARQ generalized; destination can receive blocks from multiple nodes

System Model

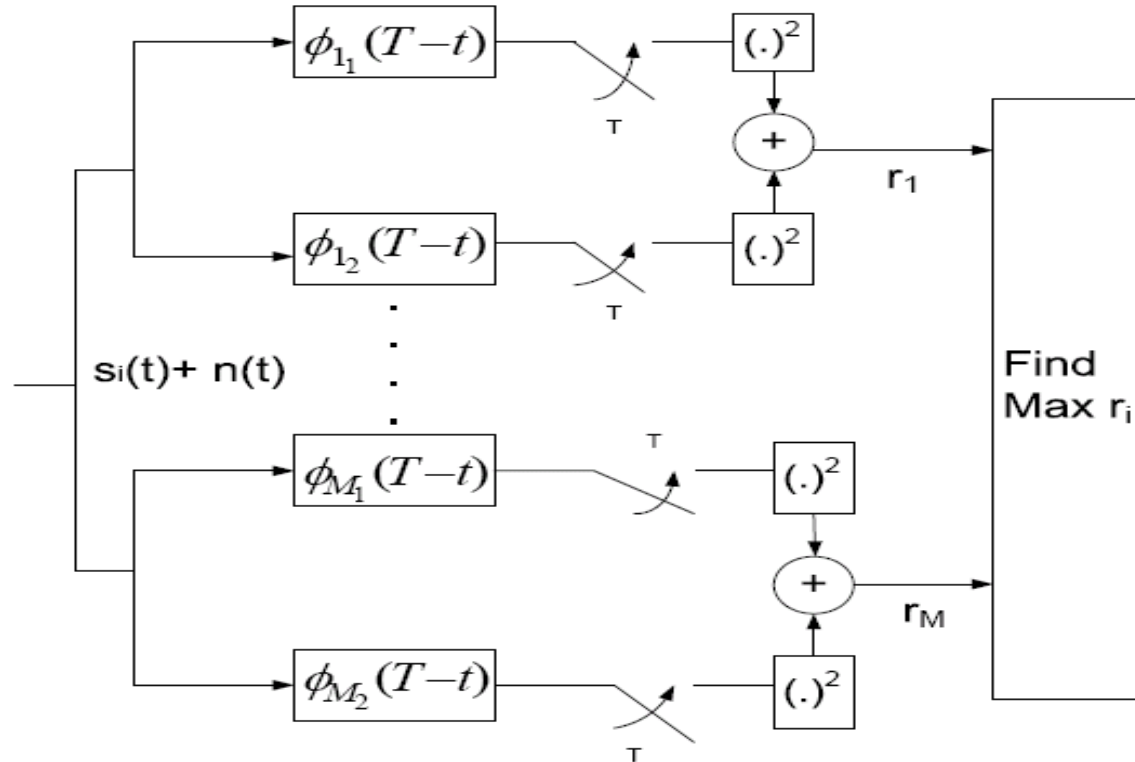


Noncoherent orthogonal modulation

- **Phase information not required**
- **M-ary modulation**
 - M basis functions; each function orthogonal to others
- **Analysis in vector space**
 - Block of size of N symbols represented by M x N matrix
 - Symbol S_n represented by n^{th} column vector; all elements zero except row m (set to one) for m^{th} basis function
 - Constant fading coefficient \mathbf{h} over a block; complex Gaussian noise added to each matrix element $\mathbf{y}[m,n]$
 - $\mathbf{y}[m,n] = \mathbf{h} \cdot \mathbf{x}[m,n] + \mathbf{v}[m,n]$

Demodulation for Noncoherent FSK

- $\phi_{i_1}(T-t)$ is match-filtered to $\cos(2\pi f_i t)$ (inphase), $\phi_{i_2}(T-t)$ corresponds to $\sin(2\pi f_i t)$ (quadrature) while f_i is the i^{th} frequency tone



Discrete Memoryless Channel (DMC) Capacity

- Mutual information

$$I(X, Y) = \sum_y \sum_x p(x, y) \log_2 \frac{p(x, y)}{p(x)p(y)} \quad C = \max_{p(x)} I(X, Y)$$

- M basis functions have zero correlation. Apriori probability of any symbol is 1/M so an error implies possibility of $M-1$ equi-probable symbols

- Symbol error probability

$$P_s = \sum_{k=1}^{M-1} \binom{M-1}{k} \log_2 \frac{(-1)^{k+1}}{k+1} \exp \left[-\frac{k}{k+1} \frac{E_s}{N_o} \right]$$

DMC Capacity

- Using P_s the $M \times M$ symbol transition matrix T_s is generated

$$T_s = \begin{bmatrix} 1 - P_s & \frac{P_s}{M-1} & \cdot & \frac{P_s}{M-1} \\ \frac{P_s}{M-1} & 1 - P_s & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \frac{P_s}{M-1} & \cdot & \cdot & 1 - P_s \end{bmatrix}$$

- $T_s(i,j) = \mathbf{p(j|i)}$ conditional probability that i^{th} symbol sent and j^{th} received
- Using T_s Mutual information (i.e. DMC capacity) is computed as

$$C = \log_2 M + P_s \log_2 \frac{P_s}{M-1} + (1 - P_s) \log_2 (1 - P_s)$$

Block Fading

■ Slow fading environment (Erdogic capacity does not exist)

- Encoded codeword broken into B blocks each with rate R_B
- SNR exponentially distributed due to Rayleigh fading
- Code-combined capacity is given by:

$$C_B = \log_2 M + \overline{P_s} \log_2 \frac{\overline{P_s}}{M-1} + (1 - \overline{P_s}) \log_2 (1 - \overline{P_s})$$

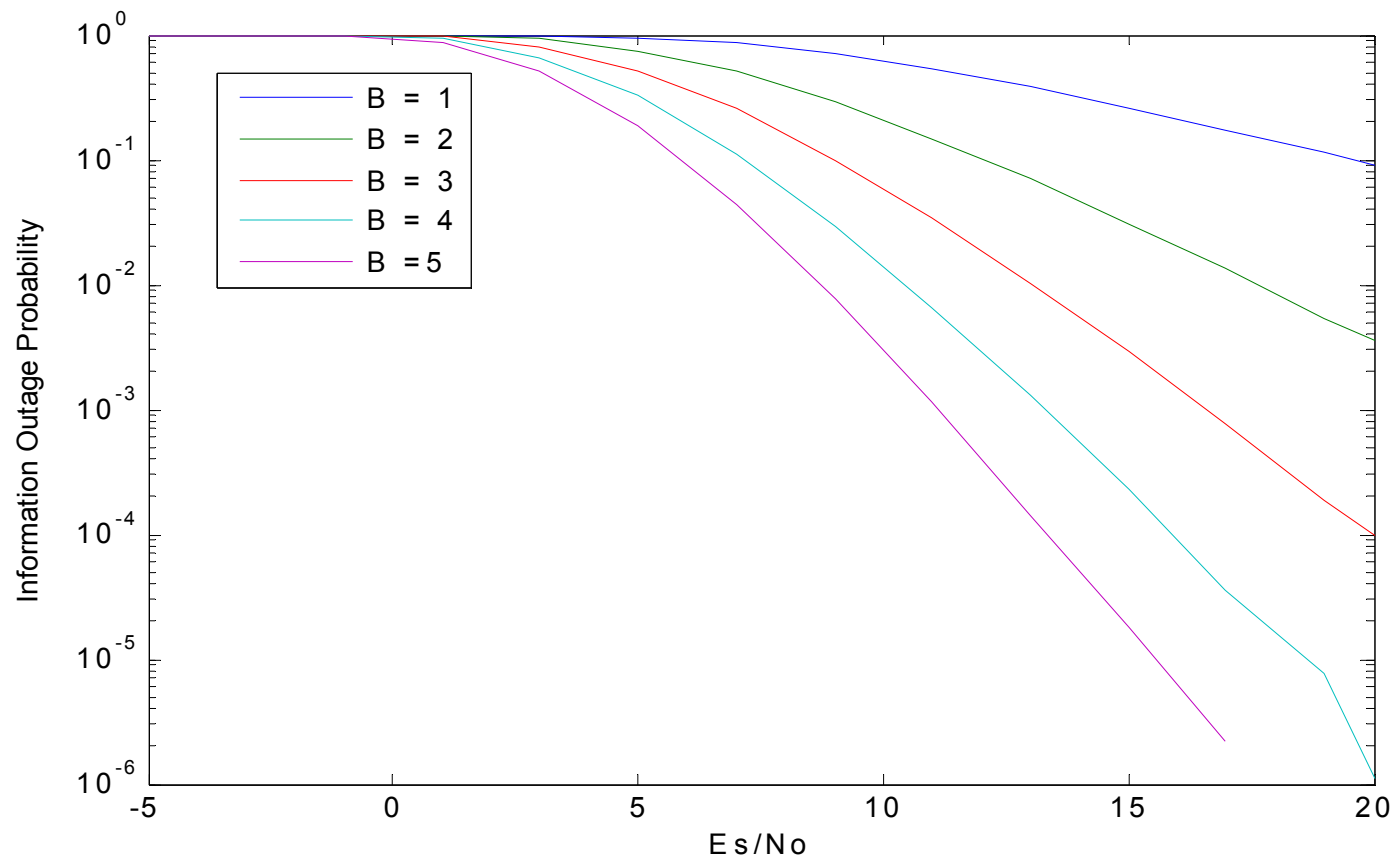
$$\text{where } \overline{P_s} = \frac{P_{s_1} + P_{s_2} + \dots + P_{s_B}}{B}$$

■ Information Outage Probability $P_0(B)$

- $P_0(B) = P[C_B < R_B]$
- For example, for uncoded 16-ary NFSK, $R_B = \log_2 16 = 4$

Code-combining in Block Fading

- Monte-Carlo simulations used. In a (51, 45) block code, rate $R_B = 4 \times (45/51)$



Hybrid-ARQ

- Combines **Forward Error Correction** (FEC) with **Automatic Repeat Request** (ARQ).
- A message of length k is encoded to a n length codeword. Codeword is broken into B_{max} blocks.
- Each of the blocks *incrementally* sent over the channel until the destination able to decode the message or all have been sent.
- Types of hybrid-ARQ
 - Type I:** Individual blocks decoded but not combined
 - Type II:** All accumulated blocks used for decoding but cannot be decoded on their own.
 - Type III:** Individual blocks also decode-able

Reed Solomon as Channel Code

■ RS (n,k)

- RS codes are nonbinary.
- can correct $t = (n-k)/2$ symbol errors
- n is number of encoded symbols
- K is number of message symbols
- For 256-ary RS symbols each symbol corresponds to one byte of data.
- k bytes are input to the encoder, where $k \leq 51$.
- Output of encoder is a 255 byte long mother code word
- Our hybrid-ARQ protocol breaks it into 5 blocks, each of length 51 bytes.
- In 16-ary modulation each RS symbol matched to two modulation symbols

Symbol Errors Decoding

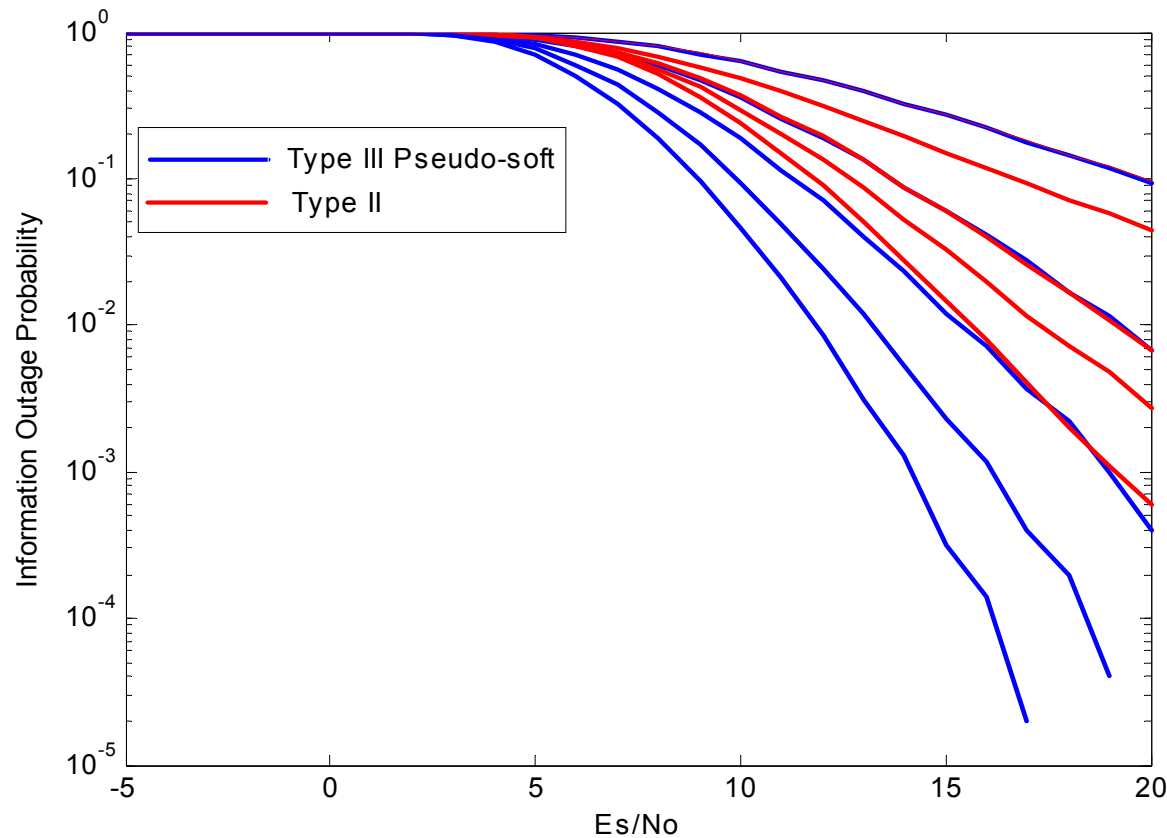
- Consider $k=45$
 - The number of correctable errors after each block transmission is:

block	n	t
1	51	3
2	102	28
3	153	54
4	204	79
5	255	105

- Performance can be poor if one block has many errors
 - Iteratively decode error blocks by trying different combinations; we can call this selective-combining technique **Pseudo-soft** decoding
- For example by $B = 4$, block errors are 14, 51, 21, 17. Since total Errors are 104 the codeword is uncorrectable
- Pseudo-soft approach decodes 53 errors for blocks 1, 3 and 4

Gains from Selective Code-Combining

- Results for a (255, 45) RS mother code using noncoherent orthogonal modulation ($M=16$)



Throughput Derivation

- Pmf of B, the number of hybrid-ARQ transmissions until successful decoding is denoted by

$$p_B[b] = \begin{cases} \xi(1 - p_o(b)) \prod_{i=1}^{b-1} p_o(i) & \text{for } 1 \leq b \leq B_{\max} \end{cases}$$

where

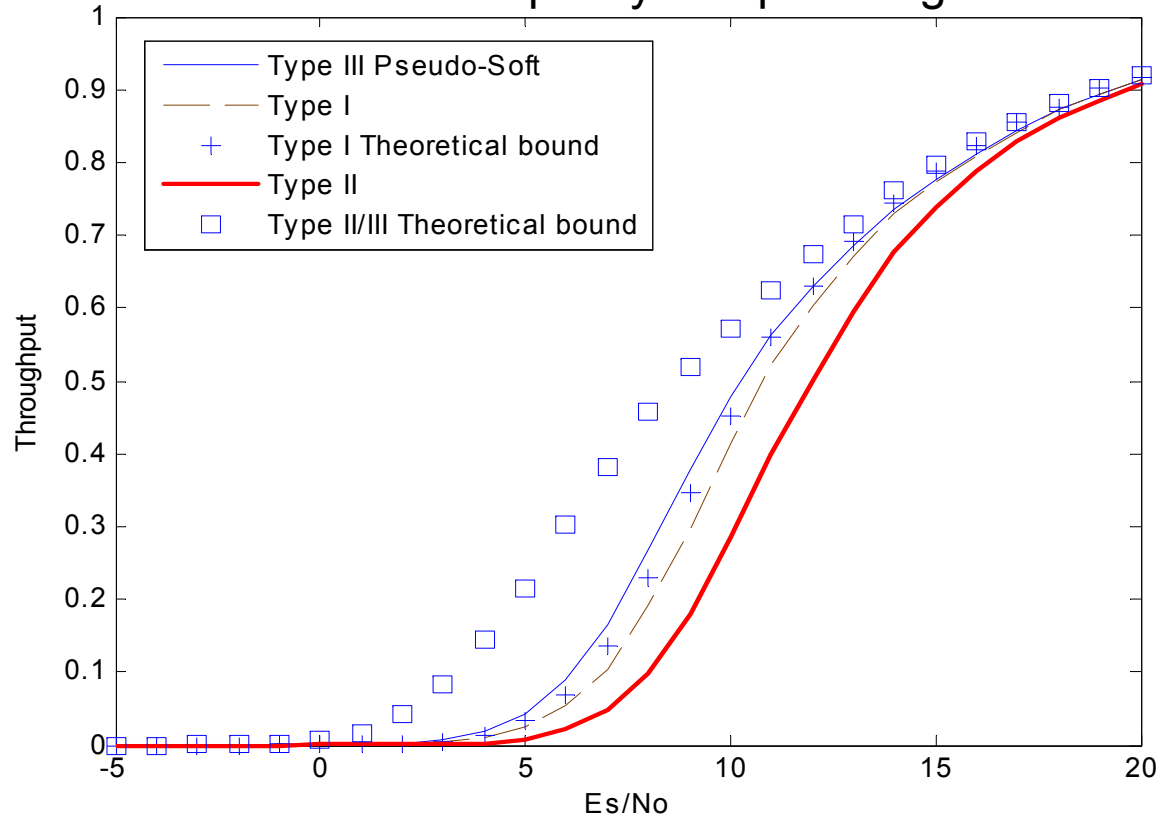
$$\xi = \left[\sum_{i=1}^{B_{\max}} (1 - p_b(b)) \prod_{i=1}^{b-1} p_b(b) \right]$$

- Throughput efficiency is the ratio of correct bits to transmitted bits

$$\eta_{eff} = \frac{1 - p_0(B_{\max})}{E[B]}$$

Throughput Performance Comparison

- Noncoherent 16-ary orthogonal modulation based hybrid-ARQ with a (255, 45) Reed Solomon mother code assumed. The *Theoretical bound* curves represent the results for DMC capacity compared against rate $R_B = (45/51) \times 4$



Network Relaying

■ Decode and Forward relaying

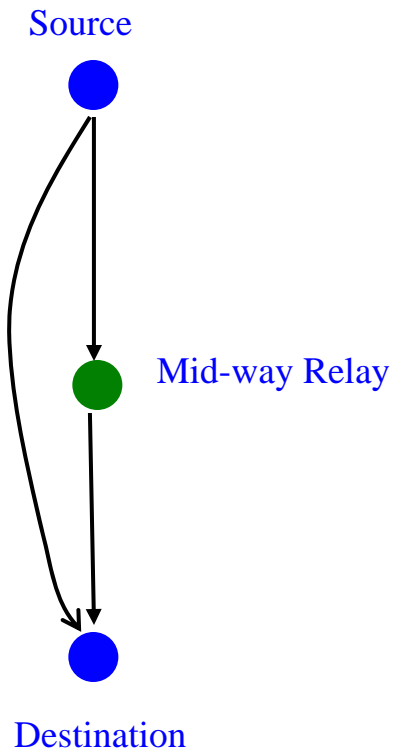
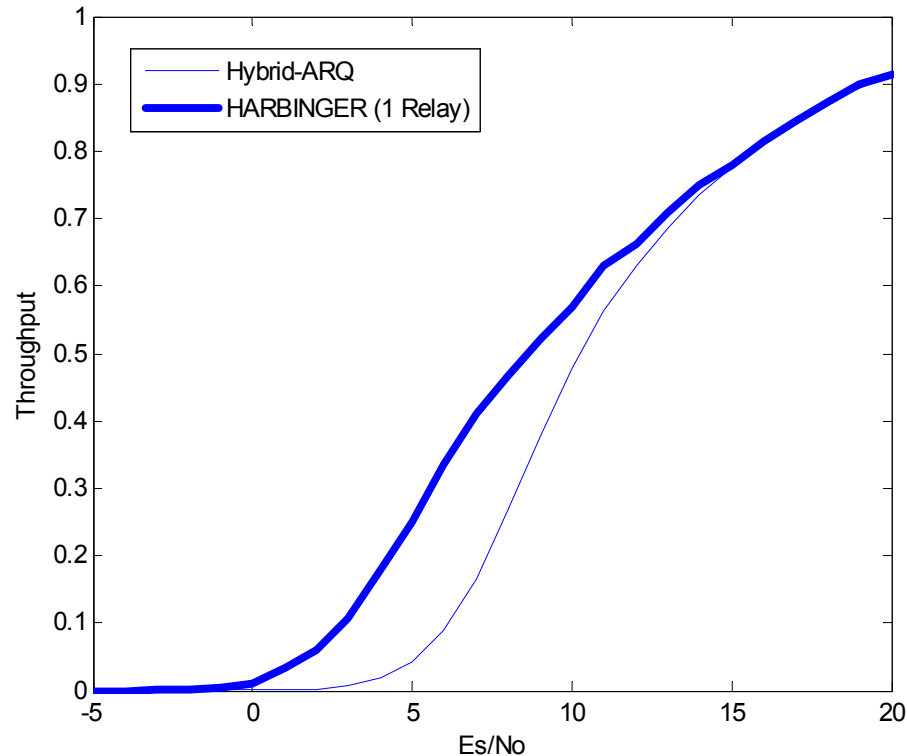
- After each ARQ transmission, some of the intermediate nodes could “overhear” the transmissions.
- Overhearing nodes that correctly decode could serve as **relays**. The ARQ retransmission could come from a relay instead of the source.

■ HARBINGER (Hybrid-ARQ based intra-cluster geographic relaying)

- can Source broadcasts first packet
- Relays that can decode are added to the **decoding set D** . The source is already in D
- Next packet sent by a node in D geographically closest to destination
- Process continues until the destination can decode or all blocks sent.

Relaying Gains

- Performance of HARBINGER with a path loss coefficient of 2. Results with modulation index $M=16$, RS codeword (255, 45) and Pseudo-soft decoding shown below.



Applications to Zigbee

■ IEEE 802.15.4/ Zigbee standard

- Pseudo-orthogonal modulation
- Offers peer-to-peer topology
- CSMA-CA mechanism for medium access control
- Assigns ID to cluster nodes

■ Programming the protocol in Applications layer

- Addresses of source, destination and the message encoded in a codeword. Its first block is transported as payload of **data frame** to lower layers
- Subsequent blocks sent in newer **data frames** in case of decoding failures

Conclusions

- A practical hybrid-ARQ protocol presented:
 - Noncoherent orthogonal modulation
 - Reed Solomon coding
 - Pseudo-soft decoding
- Information theoretic comparisons with DMC capacity more appropriate since off-the-shelf transceivers do not use matched filter outputs.
- Protocol used for relay networks
- Relevance to Zigbee-based networks shown

**Thank you
Questions?**