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Outline





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- The Binomial Theorem
 - Motivation
 - Pascal's Triangle
 - The Theorem
 - Application

Permutation

Combinations The Binomial Theorem

Permutations

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Definition

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Computing P(n, r)

$$P(n,r) =$$

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$$P(n,r) = n \cdot (n-1) \cdot \ldots (n-r+1)$$

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$$P(n,r) = n \cdot (n-1) \cdot \dots (n-r+1) \\ = n \cdot (n-1) \cdot \dots (n-r+1) \cdot \frac{(n-r) \cdot (n-r-1) \cdot \dots 1}{(n-r) \cdot (n-r-1) \cdot \dots 1}$$

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Permutations (contd.)

Example

Compute P(7,3), P(n,0), P(n,1), and P(n,n).

Permutations (contd.)

Example

Compute P(7,3), P(n,0), P(n,1), and P(n, n). Solution: 210,

Permutations (contd.)

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Compute P(7,3), P(n,0), P(n,1), and P(n,n). Solution: 210, 1,

Permutations (contd.)

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Compute *P*(7,3), *P*(*n*,0), *P*(*n*,1), and *P*(*n*,*n*). Solution: 210, 1, *n*,

Permutations (contd.)

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Compute *P*(7,3), *P*(*n*,0), *P*(*n*,1), and *P*(*n*,*n*). Solution: 210, 1, *n*, and *n*!.

Permutations (contd.)

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Compute P(7, 3), P(n, 0), P(n, 1), and P(n, n). Solution: 210, 1, *n*, and *n*!.

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How many 3 letter words can be formed using the letters in the word "compiler"?

Permutations (contd.)

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How many 3 letter words can be formed using the letters in the word "compiler"? Solution: P(8,3).

Permutations (contd.)

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How many 3 letter words can be formed using the letters in the word "compiler"? Solution: P(8,3).

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In how many ways can a president and vice-president be chosen from a group of 20 people?

Permutations (contd.)

Example

Compute P(7, 3), P(n, 0), P(n, 1), and P(n, n). Solution: 210, 1, *n*, and *n*!.

Example

How many 3 letter words can be formed using the letters in the word "compiler"? Solution: P(8,3).

Example

In how many ways can a president and vice-president be chosen from a group of 20 people? Solution: P(20, 2).

One more exampe

Example

A library has 4 books on programming, 7 on algorithms and 3 on complexity. In how many ways can the books be ordered on a shelf?

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A library has 4 books on programming, 7 on algorithms and 3 on complexity. In how many ways can the books be ordered on a shelf? Provided that the books of a subject are required to be together? Solution: If there is no restriction, the number of arrangements is P(14, 14) = 14!. Now consider the case in which the books of a given subject are required to be together. First arrange the three subjects. This can be done in P(3,3) = 3! ways. Corresponding to each such arrangement, the programming books can be permuted in P(4,4) = 4! ways, the algorithms books can be permuted in P(7,7) = 7! ways and the complexity books can be permuted in P(3,3) = 3! ways. Using the multiplication principle, the total number of arrangements is $3! \cdot 4! \cdot 7! \cdot 3!$.

Combinations

Definition

Subramani Sets and Combinatorics

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Computing C(n, r)

Focus on a given combination of r objects chosen from n objects. The objects in this combination can be permuted in r! different ways to get r! distinct permutations.

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Example

Compute *C*(7,3), *C*(*n*,0), *C*(*n*,1) and *C*(*n*,*n*). Solution: 35, 1, *n*, 1.

Combinations (examples)

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A committee of 8 students is to be selected from 19 freshmen and 34 sophomores. In how many ways, can this committee be formed, if

• it must contain 3 freshmen and 5 sophomores.

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• it must contain 3 freshmen and 5 sophomores. Solution: $C(19,3) \cdot C(34,5)$.

2 it must contain exactly one freshman.

Combinations (examples)

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- it must contain 3 freshmen and 5 sophomores. Solution: $C(19,3) \cdot C(34,5)$.
- 2 it must contain exactly one freshman. Solution: $C(19, 1) \cdot C(34, 7)$.

Combinations (examples)

Example

- it must contain 3 freshmen and 5 sophomores. Solution: $C(19,3) \cdot C(34,5)$.
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- 3 it can contain at most one freshman.

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- \bigcirc it can contain at most one freshman. Solution: C(34, 8)

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- 2 it must contain exactly one freshman. Solution: $C(19, 1) \cdot C(34, 7)$.
- 3 it can contain at most one freshman. Solution: $C(34, 8) + C(19, 1) \cdot C(34, 7)$.

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- 3 it can contain at most one freshman. Solution: $C(34, 8) + C(19, 1) \cdot C(34, 7)$.
- it contains at least one freshman. Solution: C(53, 8) C(34, 8).

Motivation Pascal's Triangle The Theorem Application

Motivation Pascal's Triangle Fhe Theorem Application

Outline



2 Combinations

The Binomial Theorem Motivation

- Pascal's Triangle
- The Theorem
- Application

Motivation Pascal's Triangle The Theorem Application

Motivation

Expansions

Subramani Sets and Combinatorics

Motivation Pascal's Triangle The Theorem Application

Motivation

Expansions

(i) $(a+b)^1 =$

Motivation Pascal's Triangle The Theorem Application

Motivation

(i)
$$(a+b)^1 = a+b$$
.

Motivation Pascal's Triangle The Theorem Application

Motivation

(i)
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.
(ii) $(a+b)^2 =$

Motivation Pascal's Triangle The Theorem Application

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(i)
$$(a+b)^1 = a+b$$
.

(ii)
$$(a+b)^2 = a^2 + 2ab + b^2$$
.

Motivation Pascal's Triangle The Theorem Application

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$$(a+b)^1 = a+b$$
.

(ii)
$$(a+b)^2 = a^2 + 2ab + b^2$$
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$$(a+b)^3 =$$

Motivation Pascal's Triangle The Theorem Application

Motivation

- (i) $(a+b)^1 = a+b$.
- (ii) $(a+b)^2 = a^2 + 2ab + b^2$.
- (iii) $(a+b)^3 = a^3 + 3a^2b + 3ab^2 + b^3$.

Motivation Pascal's Triangle The Theorem Application

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(iv) $(a + b)^4 = ???$

Motivation Pascal's Triangle The Theorem Application

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Expansions

- (i) $(a+b)^1 = a+b$.
- (ii) $(a+b)^2 = a^2 + 2ab + b^2$.
- (iii) $(a+b)^3 = a^3 + 3a^2b + 3ab^2 + b^3$.

(iv)
$$(a+b)^4 = ???$$

We want a general formula that permits us to write down the terms of $(a + b)^n$ without actual multiplication.

Permutations Combinations Motivation

Outline





The Binomial Theorem

- Motivation
- Pascal's Triangle
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Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

The coefficient table

Consider the following table:

Subramani Sets and Combinatorics

Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

The coefficient table

Consider the following table:

Row 0:

C(0, 0)

Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

The coefficient table			
Consider the following table:			
Row 0:	<i>C</i> (0, 0)	
Row 1:	<i>C</i> (1, 0)	<i>C</i> (1, 1)	

Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

The coefficient table

Consider the following table:

Row 0:			C(0, 0)		
Row 1:		C(1,0)		<i>C</i> (1, 1)	
Row 2:	C(2, 0)		C(2, 1)		C(2, 2)

Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

The coefficient table

Consider the following table:

Row 0:				C(0, 0)			
Row 1:			C(1,0)		C(1, 1)		
Row 2:		C(2, 0)		C(2, 1)		C(2, 2)	
Row 3:	C(3, 0)		C(3, 1)		C(3, 2)		C(3, 3)

Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

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Motivation Pascal's Triangle The Theorem Application

Pascal's Triangle

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Row 0:			C(0, 0)				
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Row 3:	C(3, 0)	C(3, 1)		C(3, 2)		C(3, 3)	
:							
Row <i>n</i> : <i>C</i> (<i>n</i> , 0)	C(n, 1)				<i>C</i> (<i>n</i> , <i>n</i> - 1)		C(n, n)

Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

Writing down the values of the terms gives the following table:
Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

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Row 0:

Subramani Sets and Combinatorics

1

Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

Row 0:		1	
Row 1:	1		1

Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

Row 0:			1	
Row 1:		1		1
Row 2:	1		2	

Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

Row 0:				1			
Row 1:			1		1		
Row 2:		1		2		1	
Row 3:	1		3		3		

Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

Row 0:				1			
Row 1:			1		1		
Row 2:		1		2		1	
Row 3:	1		3		3		1

Motivation Pascal's Triangle The Theorem Application

Pascal's triangle (contd.)

The Value Table

Row 0:					1				
Row 1:				1		1			
Row 2:			1		2		1		
Row 3:		1		3		3		1	
- ·									
Row n:	1		n				n		1

Motivation Pascal's Triangle The Theorem Application

Pascal's formula

Theorem

 $C(n, k) = C(n - 1, k - 1) + C(n - 1, k), 1 \le k \le n - 1.$

Motivation Pascal's Triangle Fhe Theorem Application

Pascal's formula

Theorem

 $C(n, k) = C(n - 1, k - 1) + C(n - 1, k), 1 \le k \le n - 1.$

Proof.

$$C(n-1, k-1) + C(n-1, k) = \frac{(n-1)!}{(k-1)![(n-1-(k-1))!]} + \frac{(n-1)!}{k!(n-1-k)!}$$

Motivation Pascal's Triangle The Theorem Application

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$$= \frac{(n-1)!}{(k-1)!(n-k)!} + \frac{(n-1)!}{k!(n-1-k)!}$$

Motivation Pascal's Triangle The Theorem Application

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$$= \frac{(n-1)!}{(k-1)!(n-k)!} + \frac{(n-1)!}{k!(n-1-k)!}$$
$$= \frac{k(n-1)!}{k!(n-k)!} + \frac{(n-k)(n-1)!}{k!(n-k)!}$$

Motivation Pascal's Triangle The Theorem Application

Pascal's formula

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 $C(n, k) = C(n - 1, k - 1) + C(n - 1, k), 1 \le k \le n - 1.$

Proof.

$$\begin{aligned} C(n-1,k-1)+C(n-1,k) &= \frac{(n-1)!}{(k-1)![(n-1-(k-1)!]} + \frac{(n-1)!}{k!(n-1-k)!} \\ &= \frac{(n-1)!}{(k-1)!(n-k)!} + \frac{(n-1)!}{k!(n-1-k)!} \\ &= \frac{k(n-1)!}{k!(n-k)!} + \frac{(n-k)(n-1)!}{k!(n-k)!} \\ &= \frac{(n-1)!}{k!(n-k)!} [k+(n-k)] \end{aligned}$$

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Theorem

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$$= \frac{(n-1)!}{k!(n-k)!} [k + (n-k)]$$

$$= \frac{n(n-1)!}{k!(n-k)!}$$

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Pascal's formula

Theorem

 $C(n, k) = C(n - 1, k - 1) + C(n - 1, k), 1 \le k \le n - 1.$

Proof.

$$C(n-1, k-1) + C(n-1, k) = \frac{(n-1)!}{(k-1)![(n-1-(k-1)!]} + \frac{(n-1)!}{k!(n-1-k)!}$$

$$= \frac{(n-1)!}{(k-1)!(n-k)!} + \frac{(n-1)!}{k!(n-1-k)!}$$

$$= \frac{k(n-1)!}{k!(n-k)!} + \frac{(n-k)(n-1)!}{k!(n-k)!}$$

$$= \frac{(n-1)!}{k!(n-k)!} [k + (n-k)]$$

$$= \frac{n!}{k!(n-k)!}$$

$$= \frac{n!}{k!(n-k)!}$$

Motivation Pascal's Triangle The Theorem Application

Alternative Proof

A second Proof

Observe that C(n, k) represents the number of ways in which k objects can be selected from *n* objects.

Motivation Pascal's Triangle The Theorem Application

Alternative Proof

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Observe that C(n, k) represents the number of ways in which k objects can be selected from *n* objects. Focus on a particular object, say *o*.

Motivation Pascal's Triangle The Theorem Application

Alternative Proof

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Observe that C(n, k) represents the number of ways in which k objects can be selected from n objects. Focus on a particular object, say o. Note that each selection of k objects from the n objects, either includes o or it does not.

Motivation Pascal's Triangle The Theorem Application

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Observe that C(n, k) represents the number of ways in which k objects can be selected from n objects. Focus on a particular object, say o. Note that each selection of k objects from the n objects, either includes o or it does not. Let T_1 denote the number of ways in which k objects are selected from the n objects, with o definitely included.

Motivation Pascal's Triangle The Theorem Application

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Motivation Pascal's Triangle The Theorem Application

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Motivation Pascal's Triangle The Theorem Application

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Observe that C(n, k) represents the number of ways in which k objects can be selected from n objects. Focus on a particular object, say o. Note that each selection of k objects from the n objects, either includes o or it does not. Let T_1 denote the number of ways in which k objects are selected from the n objects, with o definitely included. But this means that we have to choose (k - 1) objects from the remaining (n - 1) objects, i.e., $T_1 = C(n - 1, k - 1)$. Let T_2 denote the number of ways in which k objects are selected from the n objects, i.e., $T_2 = C(n - 1, k - 1)$.

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Motivation Pascal's Triangle The Theorem Application

Alternative Proof

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Observe that C(n, k) represents the number of ways in which k objects can be selected from n objects. Focus on a particular object, say o. Note that each selection of k objects from the n objects, either includes o or it does not. Let T_1 denote the number of ways in which k objects are selected from the n objects, with o definitely included. But this means that we have to choose (k - 1) objects from the remaining (n - 1) objects, i.e., $T_1 = C(n - 1, k - 1)$. Let T_2 denote the number of ways in which k objects are selected from the n objects, with o definitely excluded. But this means that all k objects are selected from the remaining (n - 1) objects, i.e., $T_2 = C(n - 1, k)$. Using the addition principle, $C(n, k) = T_1 + T_2 = C(n - 1, k - 1) + C(n - 1, k)$.

Note

The above proof is called a combinatorial proof and is always preferred on account of its elegance.

Motivation

Outline





The Binomial Theorem

- Motivation
- Pascal's Triangle
- The Theorem
- Application

Motivation Pascal's Triangle The Theorem Application

The Theorem

Theorem

$$(a+b)^n = \sum_{i=0}^n C(n,i) a^{n-i} \cdot b^i, \ \forall n \ge 0.$$

Motivation Pascal's Triangle The Theorem Application

The Theorem

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$$(a+b)^n = \sum_{i=0}^n C(n,i) a^{n-i} \cdot b^i, \ \forall n \ge 0.$$

Proof.

Let P(n) denote the proposition in the above theorem. We prove that P(n) holds for all n by using mathematical induction.

Motivation Pascal's Triangle The Theorem Application

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Let P(n) denote the proposition in the above theorem. We prove that P(n) holds for all n by using mathematical induction.

BASIS: At n = 0, the LHS is $(a + b)^0 = 1$

Motivation Pascal's Triangle The Theorem Application

The Theorem

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Proof.

Let P(n) denote the proposition in the above theorem. We prove that P(n) holds for all n by using mathematical induction.

BASIS: At n = 0, the LHS is $(a + b)^0 = 1$ and the RHS is $\sum_{i=0}^{0} C(0, i) a^{0-i} \cdot b^i$.

Motivation Pascal's Triangle The Theorem Application

The Theorem

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BASIS: At n = 0, the LHS is $(a + b)^0 = 1$ and the RHS is $\sum_{i=0}^0 C(0, i) a^{0-i} \cdot b^i$. Since the only value for i is also 0, the RHS is $C(0, 0) a^0 \cdot b^0 = 1$.

Motivation Pascal's Triangle The Theorem Application

The Theorem

Theorem

$$(a+b)^n = \sum_{i=0}^n C(n,i) a^{n-i} \cdot b^i, \ \forall n \ge 0.$$

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Let P(n) denote the proposition in the above theorem. We prove that P(n) holds for all n by using mathematical induction.

BASIS: At n = 0, the LHS is $(a + b)^0 = 1$ and the RHS is $\sum_{i=0}^{0} C(0, i) a^{0-i} \cdot b^i$. Since the only value for *i* is also 0, the RHS is $C(0, 0) a^0 \cdot b^0 = 1$. Thus, LHS = RHS and the basis is proven.

Motivation Pascal's Triangle The Theorem Application

Proof of Binomial Theorem

Proof.

INDUCTIVE STEP: Assume that P(k) is true, for some $k \ge 0$, i.e., assume that

$$(a+b)^{k} = \sum_{i=0}^{k} C(k,i) a^{k-i} \cdot b^{i}, \text{ for some } k \geq 0.$$

Motivation Pascal's Triangle The Theorem Application

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Motivation Pascal's Triangle The Theorem Application

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$$(a+b)^k = \sum_{i=0}^k C(k,i) a^{k-i} \cdot b^i$$
, for some $k \ge 0$.

At n = k + 1, we have,

 $LHS = (a+b)^{k+1}$

Motivation Pascal's Triangle The Theorem Application

Proof of Binomial Theorem

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LHS =
$$(a+b)^{k+1}$$

= $(a+b)^k(a+b)$

Motivation Pascal's Triangle The Theorem Application

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= $(a + b)^k (a + b)$
= $(\sum_{i=0}^k C(k, i) a^{k-i} \cdot b^i)$

Motivation Pascal's Triangle The Theorem Application

Proof of Binomial Theorem

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$$(a+b)^{k} = \sum_{i=0}^{k} C(k,i) a^{k-i} \cdot b^{i}, \text{ for some } k \ge 0.$$

$$LHS = (a+b)^{k+1}$$

= $(a+b)^k(a+b)$
= $(\sum_{i=0}^k C(k,i) a^{k-i} \cdot b^i) \cdot (a+b)$, using the inductive hypothesis
Permutations Combinations Motivation Pascal's Triangle The Theorem Application

Proof of Binomial Theorem

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$$(a + b)^{k+1}$$

= $(a + b)^{k}(a + b)$
= $(\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i}) \cdot (a + b)$, using the inductive hypothesis
= $a \cdot (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i}) + b \cdot (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i})$

Permutations Combinations Motivation Pascal's Triangle The Theorem Application

Proof of Binomial Theorem

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= $a \cdot (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i}) + b \cdot (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i})$
= $\sum_{i=0}^{k} C(k, i) a^{k+1-i} \cdot b^{i} + \sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i+1}$

Motivation Pascal's Triangle The Theorem Application

Proof of Binomial Theorem

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INDUCTIVE STEP: Assume that P(k) is true, for some $k \ge 0$, i.e., assume that

1. . .

$$(a+b)^k = \sum_{i=0}^k C(k,i) a^{k-i} \cdot b^i$$
, for some $k \ge 0$.

At n = k + 1, we have, LHS =

$$S = (a + b)^{k+1}$$

$$= (a + b)^{k}(a + b)$$

$$= (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i}) \cdot (a + b), \text{ using the inductive hypothesis}$$

$$= a \cdot (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i}) + b \cdot (\sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i})$$

$$= \sum_{i=0}^{k} C(k, i) a^{k+1-i} \cdot b^{i} + \sum_{i=0}^{k} C(k, i) a^{k-i} \cdot b^{i+1}$$

$$= C(k, 0) a^{k+1} \cdot b^{0} + \sum_{i=1}^{k} C(k, i) a^{k+1-i} \cdot b^{i} + \sum_{i=0}^{k-1} C(k, i) a^{k-i} \cdot b^{i+1} + C(k, k) a^{0} \cdot b^{k+1}$$

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

Proof.

We focus on the quantity

$$(\mathbf{F}) \sum_{i=1}^{k} C(k,i) \ a^{k+1-i} \cdot b^{i} + (\mathbf{S}) \sum_{i=0}^{k-1} C(k,i) \ a^{k-i} \cdot b^{i+1}$$
(1)

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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(F)
$$\sum_{i=1}^{k} C(k, i) a^{k+1-i} \cdot b^{i} + (\mathbf{S}) \sum_{i=0}^{k-1} C(k, i) a^{k-i} \cdot b^{i+1}$$
 (1)

Observe that the first *k* terms in **F** are $a^k \cdot b^1$, $a^{k-1} \cdot b^2$, ..., $a^1 \cdot b^k$, while the first *k* terms in **S** are also $a^k \cdot b^1$, $a^{k-1} \cdot b^2$, ..., $a^1 \cdot b^k$.

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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In other words, the terms in F and S are identical, except for the coefficents.

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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In other words, the terms in F and S are identical, except for the coefficents.

Further, all the terms can be generated using the term formula, $a^{k+1-\rho} \cdot b^{\rho}$, $1 \le \rho \le k$.

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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Further, all the terms can be generated using the term formula, $a^{k+1-p} \cdot b^p$, $1 \le p \le k$.

Observe that the coefficient of $a^{k+1-p} \cdot b^p$ is C(k, p) in **F** and C(k, p-1) in **S**. (This requires some thought!)

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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Accordingly, the coefficient of $a^{k+1-p} \cdot b^p$ in the sum (**F** + **S**) is C(k, p) + C(k, p-1), which is C(k + 1, p), using Pascal's formula.

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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Accordingly, the coefficient of $a^{k+1-p} \cdot b^p$ in the sum (**F** + **S**) is C(k, p) + C(k, p - 1), which is C(k + 1, p), using Pascal's formula. Thus, the LHS can be written as:

$$C(k, 0) a^{k+1} \cdot b^0 + \sum_{i=1}^k C(k+1, i) a^{k+1-i} \cdot b^i + C(k, k) a^0 \cdot b^{k+1}$$

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

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$$\sum_{i=1}^{k} C(k, i) a^{k+1-i} \cdot b^{i} + (\mathbf{S}) \sum_{i=0}^{k-1} C(k, i) a^{k-i} \cdot b^{i+1}$$
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= $C(k+1, 0) a^{k+1} \cdot b^{0} + \sum_{i=1}^{k} C(k+1, i) a^{k+1-i} \cdot b^{i} + C(k+1, k+1) a^{0} \cdot b^{k+1}$

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

Proof.

We focus on the quantity

(F)
$$\sum_{i=1}^{k} C(k, i) a^{k+1-i} \cdot b^{i} + (\mathbf{S}) \sum_{i=0}^{k-1} C(k, i) a^{k-i} \cdot b^{i+1}$$
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= $C(k+1, 0) a^{k+1} \cdot b^{0} + \sum_{i=1}^{k} C(k+1, i) a^{k+1-i} \cdot b^{i} + C(k+1, k+1) a^{0} \cdot b^{k+1}$

since C(k, 0) = C(k, k) = C(k + 1, 0) = C(k + 1, k + 1) = 1

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

Proof.

It follows that

$$LHS = \sum_{i=0}^{k+1} C(k+1, i) a^{k+1-i} \cdot b^{i}$$
$$= RHS$$

Motivation Pascal's Triangle The Theorem Application

Proof (contd.)

Proof.

It follows that

$$LHS = \sum_{i=0}^{k+1} C(k+1, i) a^{k+1-i} \cdot b^{i}$$
$$= RHS$$

We have thus shown that $P(k) \rightarrow P(k+1)$ and hence by applying the first principle of mathematical induction, we can conclude that P(n) is true, for all $n \ge 0$.

Motivation

Outline





The Binomial Theorem

- Motivation
- Pascal's Triangle
- The Theorem
- Application

Motivation Pascal's Triangle The Theorem Application

Application

Example

Expand $(x - 3)^4$.

Motivation Pascal's Triangle The Theorem Application

Application

Example

Expand $(x - 3)^4$. Solution:

$$\begin{aligned} (x-3)^4 &= C(4,0)x^4 \cdot (-3)^0 + C(4,1)x^3 \cdot (-3)^1 + C(4,2)x^2 \cdot (-3)^2 \\ &+ C(4,3)x^1 \cdot (-3)^3 + C(4,4)x^0 \cdot (-3)^4 \end{aligned}$$

Motivation Pascal's Triangle The Theorem Application

Application

Example

Expand $(x - 3)^4$. Solution:

$$(x-3)^4 = C(4,0)x^4 \cdot (-3)^0 + C(4,1)x^3 \cdot (-3)^1 + C(4,2)x^2 \cdot (-3)^2 + C(4,3)x^1 \cdot (-3)^3 + C(4,4)x^0 \cdot (-3)^4 = x^4 + 4x^3 \cdot (-3) + 6x^2 \cdot (9) + 4x \cdot (-27) + 81$$

Motivation Pascal's Triangle The Theorem Application

Application

Example

Expand $(x - 3)^4$. Solution:

$$(x-3)^4 = C(4,0)x^4 \cdot (-3)^0 + C(4,1)x^3 \cdot (-3)^1 + C(4,2)x^2 \cdot (-3)^2 + C(4,3)x^1 \cdot (-3)^3 + C(4,4)x^0 \cdot (-3)^4 = x^4 + 4x^3 \cdot (-3) + 6x^2 \cdot (9) + 4x \cdot (-27) + 81 = x^4 - 12x^3 + 54x^2 - 108x + 81$$

Motivation Pascal's Triangle The Theorem Application

One more example

Example

Show that

$$\sum_{i=0}^{n} C(n, i) = 2^{n}$$

Subramani Sets and Combinatorics

Motivation Pascal's Triangle The Theorem Application

One more example

Example

Show that

$$\sum_{i=0}^{n} C(n, i) = 2^{n}$$

Proof using the binomial theorem

As per the binomial theorem,

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Motivation Pascal's Triangle The Theorem Application

One more example

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Show that

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$$(1+1)^n = C(n, 0) \cdot (1) + C(n, 1) \cdot (1) + \dots + C(n, n) \cdot (1)$$

$$\Rightarrow \sum_{i=0}^n C(n, i) = 2^n$$

Motivation Pascal's Triangle The Theorem Application

An alternate proof

Motivation Pascal's Triangle The Theorem Application

An alternate proof

Proof using combinatorial arguments

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An alternate proof

Proof using combinatorial arguments

Consider a set S having n elements.

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An alternate proof

Proof using combinatorial arguments

Consider a set *S* having *n* elements. C(n, i) represents the number of ways in which *i* elements can be selected from *n* elements,

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Consider a set *S* having *n* elements. C(n, i) represents the number of ways in which *i* elements can be selected from *n* elements, i.e., C(n, i) represents the number of subsets of *S*, which have cardinality *i*.

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Motivation Pascal's Triangle The Theorem Application

A third proof

Proof using induction

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Motivation Pascal's Triangle The Theorem Application

A third proof

Proof using induction

BASIS: At n = 0,
Notivation Pascal's Triangle The Theorem Application

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Proof using induction

$$HS = \sum_{i=0}^{0} C(0, 0)$$

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Motivation Pascal's Triangle The Theorem Application

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Motivation Pascal's Triangle The Theorem Application

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Since LHS=RHS, the basis is proven.

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for some $k \ge 0$.

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Inductive proof (contd.)

Completing the induction

Motivation Pascal's Triangle The Theorem Application

Inductive proof (contd.)

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Completing the induction

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Motivation Pascal's Triangle The Theorem Application

Inductive proof (contd.)

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$$HS = \sum_{i=0}^{k+1} C(k+1, i)$$

= $C(k+1, 0) + \sum_{i=1}^{k} C(k+1, i) + C(k+1, k+1)$

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Inductive proof (contd.)

Completing the induction

Lŀ

$$\begin{split} & \mathcal{IS} = \sum_{i=0}^{k+1} C(k+1,i) \\ & = C(k+1,0) + \sum_{i=1}^{k} C(k+1,i) + C(k+1,k+1) \\ & = 1 + \sum_{i=1}^{k} [C(k,i) + C(k,i-1)] + 1, \; \text{Pascal's formula} \end{split}$$

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Inductive proof (contd.)

Completing the induction

Lŀ

$$dS = \sum_{i=0}^{k+1} C(k+1, i)$$

= $C(k+1, 0) + \sum_{i=1}^{k} C(k+1, i) + C(k+1, k+1)$
= $1 + \sum_{i=1}^{k} [C(k, i) + C(k, i-1)] + 1$, Pascal's formula
= $(1 + \sum_{i=1}^{k} C(k, i)) + (\sum_{i=1}^{k} C(k, i-1) + 1))$

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Inductive proof (contd.)

The last steps

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Inductive proof (contd.)

The last steps

$$= \sum_{i=0}^{k} C(k, i) + \sum_{j=0}^{k} C(k, j)$$

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Inductive proof (contd.)

The last steps

$$= \sum_{i=0}^{k} C(k, i) + \sum_{j=0}^{k} C(k, j)$$
$$= 2 \cdot \sum_{i=0}^{k} C(k, i)$$

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Inductive proof (contd.)

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$$= \sum_{i=0}^{k} C(k, i) + \sum_{j=0}^{k} C(k, j)$$
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Inductive proof (contd.)

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Thus LHS=RHS and the inductive step is proven.

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Thus LHS=RHS and the inductive step is proven. Applying the first principle of mathematical induction, we conclude that the conjecture is true.