## Summation - Tricks of the trade

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## 1 Factoids

Lemma: 1.1 Linearity of sums

$$\sum_{k=1}^{n} (c.a_k + b_k) = c. \sum_{k=1}^{n} a_k + \sum_{k=1}^{n} b_k$$
 (1)

Lemma: 1.2 Given the arithmetic progression, of the form

$$a_1, a_1 + d, a_1 + 2.d, \ldots,$$

the n<sup>th</sup> term in the sequence is given by

$$T_n = a_1 + (n-1).d (2)$$

and the sum of the first n terms is given by:

$$S_n = \frac{n}{2} \cdot (a_1 + T_n) \tag{3}$$

Lemma: 1.3 Given a geometric series of the form

$$1, x, x^2, \ldots,$$

the nth term in the sequence is given by

$$T_n = x^{n-1} (4)$$

and the sum of the first n terms is given by:

$$S_n = 1. \frac{x^n - 1}{x - 1}. (5)$$

If x < 1, we can rewrite  $S_n$  as:

$$S_n = \frac{1}{1-x}. (6)$$

**Remark:** 1.1 These formulae are slightly different from the ones in the text. Ensure that you understand the differences.

Remark: 1.2 Prove all the above formulae using mathematical induction.

## 2 Techniques

The following techniques are used to derive upper and lower bounds on series:

1. Telescoping - If the summand can be expressed as the difference between adjacent terms of the corresponding sequence, then telescoping can be used. Clearly, in sequence

$$a_0, a_1, \ldots, a_n$$

$$\sum_{k=1}^{n} a_k - a_{k-1} = a_n - a_0$$

Similarly,

$$\sum_{k=0}^{n-1} a_k - a_{k+1} = a_0 - a_n$$

Let us say that we want

$$\sum_{k=1}^{n-1} \frac{1}{k \cdot (k+1)}$$

Observe that  $\frac{1}{k \cdot (k+1)}$  can be rewritten as:

$$\frac{1}{k} - \frac{1}{k+1}$$

Applying the telescoping formula,

$$\sum_{k=1}^{n-1} \frac{1}{k \cdot (k+1)} = 1 - \frac{1}{n}$$

Remark: 2.1 Make sure that you understand every step!

2. Mathematical Induction - The hammer! It can be used almost everywhere. For deriving bounds, it is necessary to guess a value and then validate the guess. (Unfortunately, there will not be a samaritan lurking in the exam/quiz to give you the answer. You will have to make an intelligent guess and then check it.) You do not have to guess the exact function e.g.  $\frac{n \cdot (n+1)}{2}$ ; it suffices to guess the form i.e.  $O(n^2)$ . Consider the geometric series  $\sum_{k=0}^{n} 3^k$ . You can apply the formula, which gives  $\frac{1}{2} \cdot (3^{n+1} - 1)$  i.e.  $\frac{3}{2} \cdot 3^n - \frac{1}{2}$  i.e.  $O(3^n)$ . Or you could guess the form as  $O(3^n)$ . This means that we have to prove

$$\sum_{k=0}^{n} 3^k \le c.3^n$$

for some c > 0.

Base confirmation: Putting n=0, gives  $1 \le c$  or  $c \ge 1$ . Inductive step: Assuming that the proposition holds for n=p, we get

$$\sum_{k=0}^{p} 3^k \le c.3^p$$

for some c > 0. Now,

$$\sum_{k=0}^{p+1} 3^k = \sum_{k=0}^p 3^p + 3^{p+1}$$

$$\Rightarrow \sum_{k=0}^{p+1} 3^k \le c.3^p + 3^{p+1}$$

by the inductive hypothesis

$$\Rightarrow \sum_{k=0}^{p+1} 3^k \le (\frac{1}{3} + \frac{1}{c}).c.3^{p+1}$$

, which holds as long as  $(\frac{1}{3} + \frac{1}{c}) \le 1$  i.e.  $c \ge \frac{3}{2}$ . Thus we have proved our case i.e.

$$\sum_{k=0}^{n} 3^k \le c.3^n$$

for some  $c \geq \frac{3}{2}$  or

$$\sum_{k=0}^{n} 3^k = O(3^n)$$

3. Term bounding - Yet another idea is to bound the individual terms of the series, e.g.

$$\sum_{k=1}^{n} k \le \sum_{k=1}^{n} n = n^2$$

For a geometric series,  $\sum_{k=0}^{n} a_k$ , if  $\frac{a_{k+1}}{a_k} \leq r < 1$ , where r is a constant, then  $a_k \leq a_0 \cdot r^k$ . Hence,

$$\sum_{k=0}^{n} a_k \le \sum_{k=0}^{\infty} a_0 r^k$$

$$= a_0 \cdot \sum_{k=0}^{\infty} r^k$$

$$= a_0 \cdot \frac{1}{1 - r}$$

**Remark:** 2.2 Use the term bounding idea to derive an upper bound for  $\sum_{k=1}^{\infty} \frac{k}{3^k}$ . (Hint: Given in textbook)

4. Sum splitting - This technique is particularly useful for deriving lower bounds e.g.

$$\sum_{k=1}^{n} k = \sum_{k=1}^{\frac{n}{2}} k + \sum_{k=\frac{n}{2}+1}^{n} k$$

$$\Rightarrow \geq 0 + \sum_{k=\frac{n}{2}+1}^{n} \frac{n}{2}$$

$$\Rightarrow \geq \left(\frac{n}{2}\right)^{2}$$

5. Integration - This is a versatile technique used in deriving upper and lower bounds when the function (summand) is not amenable to the tricks above. For a monotonically increasing function, we have

$$\int_{m-1}^{n} f(x).dx \le \sum_{k=m}^{n} f(k) \le \int_{m}^{n+1} f(x).dx$$

For a monotonically decreasing function, we have

$$\int_{m-1}^{n} f(x).dx \ge \sum_{k=m}^{n} f(k) \ge \int_{m}^{n+1} f(x).dx$$

Take  $f(k) = \frac{1}{k}$ , a monotonically decreasing function. Hence,

$$\sum_{k=1}^{n} \frac{1}{k} \ge \int_{1}^{n+1} \frac{dx}{x} = \ln(n+1)$$

Likewise,

$$\sum_{k=2}^{n} \frac{1}{k} \le \int_{1}^{n} \frac{dx}{x} = \ln n$$

and thus

$$\ln(n+1) \le \sum_{k=1}^{n} \frac{1}{x} \le 1 + \ln n$$