

Analysis of Algorithms - Scrimmage I (Solutions)

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Please attempt as many problems as you can in class. The scrimmage will not be graded, i.e. there are no points. The solutions are posted at:

<http://www.csee.wvu.edu/~ksmani/courses/fa02/cs320/cs320.html>

1. Prove using mathematical induction:

(a) $\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{n \cdot (n+1)} = \frac{n}{n+1}$.

Proof: *Base case* $P(1)$:

$$\begin{aligned} LHS &= \frac{1}{1 \cdot 2} \\ &= \frac{1}{2} \\ RHS &= \frac{1}{1+1} \\ &= \frac{1}{2} \end{aligned}$$

Thus, $LHS = RHS$ and $P(1)$ is true.

Let us assume that $P(k)$ is true, i.e.

$$\frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{k \cdot (k+1)} = \frac{k}{k+1}$$

We need to show that $P(k+1)$ is true.

$$\begin{aligned} LHS &= \frac{1}{1 \cdot 2} + \frac{1}{2 \cdot 3} + \frac{1}{3 \cdot 4} + \dots + \frac{1}{(k+1) \cdot (k+2)} \\ &= \frac{k}{k+1} + \frac{1}{(k+1) \cdot (k+2)} \quad (\text{using the inductive hypothesis}) \\ &= \frac{1}{k+1} \cdot \left(k + \frac{1}{k+2}\right) \\ &= \frac{1}{k+1} \cdot \left(\frac{k^2 + 2k + 1}{k+2}\right) \\ &= \frac{1}{k+1} \cdot \left(\frac{(k+1)^2}{k+2}\right) \\ &= \frac{k+1}{k+2} \\ &= RHS \end{aligned}$$

Thus, we have shown that $P(k) \rightarrow P(k+1)$; applying the principle of mathematical induction, we conclude that the conjecture is true. \square

- (b) $7^n - 2^n$ is divisible by 5

Proof: Base case $P(1)$: Observe that $LHS = 7^1 - 2^1 = 7 - 2 = 5$, which is divisible by 5; therefore $P(1)$ is true.

Let us assume that $P(k)$ is true for some integer k , i.e. $7^k - 2^k$ is divisible by 5 for some integer k . This means $7^k - 2^k = 5 \cdot m$ for some integer m . Now consider $P(k+1)$.

$$\begin{aligned} LHS &= 7^{k+1} - 2^{k+1} \\ &= 7 \cdot 7^k - 2 \cdot 2^k \\ &= 7 \cdot (5m + 2^k) - 2 \cdot 2^k \quad (\text{using the inductive hypothesis}) \\ &= 35m + 7 \cdot 2^k - 2 \cdot 2^k \\ &= 35m + 5 \cdot 2^k \\ &= 5 \cdot (7m + 2^k) \\ &= 5q \text{ for some } q \end{aligned}$$

Thus, the LHS is divisible by 5 and $P(k+1)$ is true. We have shown that $P(k) \rightarrow P(k+1)$; applying the principle of mathematical induction, we conclude that the conjecture is true.

\square

- (c) Show that $13^n - 6^n$ is divisible by 7.

Proof: Base Case: When $n = 1$, $13^n - 6^n = 13 - 6 = 7$ is divisible by 7; thus $P(1)$ is true.

Let us assume that $P(k)$ is true, i.e. $13^k - 6^k$ is divisible by 7, for some $k > 1$. Thus, $13^k - 6^k = 7 \cdot m$, for some integer m . Now consider $P(k+1)$.

$$\begin{aligned} 13^{k+1} - 6^{k+1} &= 13 \cdot 13^k - 6 \cdot 6^k \\ &= 13 \cdot [6^k + 7m] - 6 \cdot 6^k \\ &= 13 \cdot 7m + 13 \cdot 6^k - 6 \cdot 6^k \\ &= 13 \cdot 7m + 6^k \cdot [13 - 6] \\ &= 13 \cdot 7m + 6^k \cdot 7 \\ &= 7 \cdot [13m + 6^k] \\ &= 7 \cdot q \text{ for some integer } q \end{aligned}$$

Thus, applying the principle of mathematical induction, we conclude that the conjecture is true.

\square

- (d) Show that

$$(\cos \theta + i \sin \theta)^n = \cos n\theta + i \sin n\theta \quad \forall n \geq 1$$

Proof: Base case $P(1)$:

$$\begin{aligned} LHS &= (\cos \theta + i \sin \theta)^1 \\ &= \cos \theta + i \sin \theta \\ &= RHS \end{aligned}$$

Thus, $P(1)$ is true. Let us assume that $P(k)$ is true for some $k > 1$, i.e.

$$(\cos \theta + i \sin \theta)^k = \cos k\theta + i \sin k\theta \quad (1)$$

Let us show that $P(k+1)$ is true.

$$\begin{aligned} (\cos \theta + i \sin \theta)^{k+1} &= (\cos \theta + i \sin \theta) \cdot (\cos \theta + i \sin \theta)^k \\ &= (\cos \theta + i \sin \theta) \cdot (\cos k\theta + i \sin k\theta) \text{ using the inductive hypothesis} \\ &= \cos \theta \cdot \cos k\theta + i \cos \theta \cdot \sin k\theta + i \sin \theta \cdot \cos k\theta + i^2 \sin \theta \sin k\theta \\ &= \cos \theta \cdot \cos k\theta - \sin \theta \sin k\theta + i \cdot (\cos \theta \cdot \sin k\theta + \sin \theta \cdot \cos k\theta) \\ &= \cos(1+k)\theta + i \cdot \sin(1+k)\theta \\ &= \cos(k+1)\theta + i \cdot \sin(k+1)\theta \\ &= RHS \end{aligned}$$

We apply the principle of mathematical induction and conclude that the conjecture is true. \square

2. Compare $f(n)$ and $g(n)$ using asymptotic notation; you may either describe $f(n)$ in terms of $g(n)$ (for instance, $f(n) = O(g(n))$) or $g(n)$ in terms of $f(n)$ (for instance, $g(n) = \omega(f(n))$). Make sure that your description is as precise as possible.

(a) $f(n) = n \log^5 n$, $g(n) = n^2$

Observe that

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} &= \lim_{n \rightarrow \infty} \frac{n \log^5 n}{n^2} \\ &= \lim_{n \rightarrow \infty} \frac{\log^5 n}{n} \\ &= \lim_{n \rightarrow \infty} \frac{(5 \log^4 n) \frac{1}{n}}{1} \text{ (by applying L'Hospital's rule!)} \\ &= \lim_{n \rightarrow \infty} \frac{5 \log^4 n}{n} \\ &= \lim_{n \rightarrow \infty} \frac{5 \cdot 4 \cdot 3 \cdot 2 \cdot 1}{n} \text{ (by repeated application of L'Hospital's rule!)} \\ &= 0 \end{aligned}$$

It follows that $f(n) = o(g(n))$.

(b) $f(n) = n \log_4 n$, $g(n) = n \log_{35} n$.

Note that $f(n)$ can be written as $n \log_2 n \cdot (\frac{1}{\log_2 4})$, by using the Logarithm rules. Likewise, $g(n)$ can be written as: $n \log_2 n \cdot (\frac{1}{\log_2 35})$. It is clear that $f(n)$ and $g(n)$ differ only by a constant in their rates of growth and hence $f(n) = \Theta(g(n))$.

(c) $f(n) = \log^3 n$, $g(n) = n^{\frac{1}{3}}$.

By repeated application of L'Hospital's rule, we see that

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0$$

Thus, $f(n) = o(g(n))$.

(d) $f(n) = 2^n$, $g(n) = 2^{n+1}$

Observe that $f(n) = \frac{1}{2}g(n)$ and $g(n) = 2 \cdot f(n)$, i.e. the two functions differ at most by a constant in their rates of growth. It follows that $f(n) = \Theta(g(n))$.