Advanced Analysis of Algorithms - Scrimmage I

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1 **Summation**

- 1. Assume that 0 < |x| < 1. Derive closed forms for the following sums:
 - (a) $\sum_{k=0}^{\infty} x^k$. Solution: Observe that: $\sum_{k=0}^{\infty} x^k = 1 + x + x^2 + x^3 + \ldots = S$. Multiplying S by x we get: $x \cdot S = x(1 + x + x^2 + x^3 + \ldots)$ $= x + x^2 + x^3 + x^4 + \dots$

Observe that $S - x \cdot S = S(1 - x) = 1$ and by solving for S, we get $S = \frac{1}{1 - x}$. Thus, $\sum_{k=0}^{\infty} x^k = \frac{1}{1 - x}$. (b) $\sum_{k=0}^{\infty} k \cdot x^k$.

Solution: Using the result of problem 1(a), we have: $\sum_{k=0}^{\infty} x^k = \frac{1}{1-x}$. Differentiating both sides of the above equation with respect to x and making use of the fact that the derivative of a sum is equal to the sum of derivatives, we get:

$$\begin{aligned} &(\sum_{k=0}^{\infty} x^k)' &= (\frac{1}{1-x})' \\ &\sum_{k=0}^{\infty} k \cdot x^{k-1} &= \frac{0 \cdot (1-x) - (1) \cdot (-1)}{(1-x)^2} \\ &\sum_{k=0}^{\infty} k \cdot x^{k-1} &= \frac{0+1}{(1-x)^2} \\ &\sum_{k=0}^{\infty} k \cdot x^{k-1} &= \frac{1}{(1-x)^2} \end{aligned}$$

Multiplying both sides of the above equation by x, we get $\sum_{k=0}^{\infty} k \cdot x^k = \frac{x}{(1-x)^2}$.

(c) $\sum_{k=0}^{\infty} k^2 \cdot x^k$. **Solution:** Using the result of problem 1(b), we have: $\sum_{k=0}^{\infty} k \cdot x^k = \frac{x}{(1-x)^2}$. Differentiating both sides of the above equation with respect to x and making use of the fact that the derivative of a sum is equal to the sum of derivatives, we get:

$$\begin{aligned} &(\sum_{k=0}^{\infty} k \cdot x^k)' &= (\frac{x}{(1-x)^2})' \\ &\sum_{k=0}^{\infty} k^2 \cdot x^{k-1} &= \frac{1 \cdot (1-x)^2 - x \cdot (2 \cdot (1-x) \cdot (-1))}{(1-x)^4} \end{aligned}$$

$$\begin{split} \sum_{k=0}^{\infty} k^2 \cdot x^{k-1} &= \frac{(1-x)^2 + 2 \cdot x \cdot (1-x)}{(1-x)^4} \\ \sum_{k=0}^{\infty} k^2 \cdot x^{k-1} &= \frac{(1-x)[1-x+2 \cdot x]}{(1-x)^4} \\ \sum_{k=0}^{\infty} k^2 \cdot x^{k-1} &= \frac{1+x}{(1-x)^3} \end{split}$$

Multiplying both sides of the above equation by x, we get: $\sum_{k=0}^{\infty} k^2 \cdot x^k = \frac{x(1+x)}{(1-x)^3}$. \Box

2. Show that $\sum_{k=1}^{n} \frac{1}{2 \cdot k - 1} = \ln(\sqrt{n}) + O(1)$. **Solution:** Observe that $\sum_{k=1}^{n} \frac{1}{2 \cdot k - 1} = 1 + \frac{1}{3} + \frac{1}{5} + \dots$ and that $\sum_{k=1}^{n} \frac{1}{k} = 1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots$ By subtracting $\frac{1}{2} \cdot \sum_{k=1}^{n} \frac{1}{k}$ from $\sum_{k=1}^{n} \frac{1}{k}$ we get:

$$\sum_{k=1}^{n} \frac{1}{k} - \frac{1}{2} \cdot \sum_{k=1}^{n} \frac{1}{k} = 1 + \frac{1}{3} + \frac{1}{5} + \dots$$
$$= \sum_{k=1}^{n} \frac{1}{2 \cdot k - 1}$$

Using the fact that $H_n = \sum_{i=1}^n \frac{1}{i} = \ln(n) + O(1)$, we have:

$$\sum_{k=1}^{n} \frac{1}{2 \cdot k - 1} = H_n - \frac{1}{2} \cdot H_n$$

= $ln(n) + O(1) - \frac{1}{2}(ln(n) + O(1))$
= $\frac{1}{2} \cdot ln(n) + O(1)$
= $ln(\sqrt{n}) + O(1)$

Thus, $\sum_{k=1}^{n} \frac{1}{2k-1} = ln(\sqrt{n}) + O(1)$. \Box

3. Show that $\sum_{k=0}^{\infty} \frac{k-1}{2^k} = 0$. Solution: Observe that we can rewrite the sum as,

$$\sum_{k=0}^{\infty} \frac{k-1}{2^k} = \sum_{k=0}^{\infty} \left(\frac{k}{2^k} - \frac{1}{2^k}\right)$$
$$= \sum_{k=0}^{\infty} \frac{k}{2^k} - \sum_{k=0}^{\infty} \frac{1}{2^k}$$

Notice that $\sum_{k=0}^{\infty} \frac{k}{2^k} = \sum_{k=0}^{\infty} k \cdot (\frac{1}{2})^k$. Now, we can use the result obtained from problem 1(b) in order to get

$$\sum_{k=0}^{\infty} k \cdot \left(\frac{1}{2}\right)^k = \frac{\frac{1}{2}}{(1-\frac{1}{2})^2} \\ = \frac{\frac{1}{2}}{(\frac{1}{2})^2} \\ = \frac{\frac{1}{2}}{\frac{1}{4}}$$

$$= \frac{4}{2}$$
$$= 2$$

Next, notice that $\sum_{k=0}^{\infty} \frac{1}{2^k} = \sum_{k=0}^{\infty} (\frac{1}{2})^k$, which is an infinite decreasing geometric series. Thus, $\sum_{k=0}^{\infty} (\frac{1}{2})^k = \frac{1}{1-\frac{1}{2}} = 2$.

Now, by substituting back into our equation, we get

$$\sum_{k=0}^{\infty} \frac{k}{2^k} - \sum_{k=0}^{\infty} \frac{1}{2^k} = 2 - 2$$

= 0

2 Counting

Let $C(n,k), k \leq n$, denote the number of ways of selecting k objects from a set of n objects.

1. Show that C(n,k) = C(n-1,k) + C(n-1,k-1). Solution: Observe that

$$\begin{aligned} C(n-1,k) + C(n-1,k-1) &= \frac{(n-1)!}{k! \cdot (n-1-k)!} + \frac{(n-1)!}{(k-1)! \cdot (n-k)!} \\ &= \frac{(n-1)!}{(k-1)! \cdot k \cdot (n-1-k)!} + \frac{(n-1)!}{(k-1)! \cdot (n-1-k)! \cdot (n-k)} \\ &= \frac{(n-1)! \cdot (n-k) + (n-1)! \cdot k}{(k-1)! \cdot k \cdot (n-1-k)! \cdot (n-k)} \\ &= \frac{(n-1)! \cdot n}{(k-1)! \cdot k \cdot (n-1-k)! \cdot (n-k)} \\ &= \frac{n!}{k! \cdot (n-k)!} \end{aligned}$$

Thus, $C(n,k) = C(n-1,k) + C(n-1,k-1) = \frac{n!}{k! \cdot (n-k)!}$.

2. Show that $C(n,k) = \frac{n}{k}C(n-1,k-1)$. Solution: Observe that

$$\frac{n}{k} \cdot C(n-1, k-1) = \frac{n}{k} \cdot \frac{(n-1)!}{(k-1)! \cdot (n-k)!} \\ = \frac{n \cdot (n-1)!}{k \cdot (k-1)! \cdot (n-k)!} \\ = \frac{n!}{k! \cdot (n-k)!}$$

Thus, $C(n,k) = \frac{n}{k} \cdot C(n-1,k-1)$. \Box

3. Prove that $\sum_{i=1}^{n} i = C(n+1,2)$. **Proof:** Observe that $\sum_{i=1}^{n} i = 1+2+3+\ldots = \frac{n \cdot (n+1)}{2}$. We also have that

$$C(n+1,2) = \frac{(n+1)!}{2! \cdot (n-1)!} = \frac{(n+1)!}{2 \cdot (n-1)!}$$

$$= \frac{(n+1) \cdot n \cdot (n-1)!}{2 \cdot (n-1)!}$$
$$= \frac{(n+1) \cdot n}{2}$$

Thus, $\sum_{i=1}^{n} i = C(n+1,2)$. \Box

3 Probability

Let X and Y random variables defined on a sample space S.

1. Show that $\mathbf{E}[a] = a$, if a is a constant. Solution: Let g(x) = a for all x. Then

$$\mathbf{E}[a] = \mathbf{E}[g(x)]$$

$$= \sum_{x \cdot p(x) > 0} g(x) \cdot p(x)$$

$$= \sum_{x \cdot p(x) > 0} a \cdot p(x)$$

$$= a \cdot \sum_{x \cdot p(x) > 0} p(x)$$

$$= a \cdot 1$$

$$= a$$

Thus, the expectation of a constant is the constant. \Box

- 2. If X and Y are non-negative, show that $\mathbf{E}[\max(X, Y)] = \mathbf{E}[X] + \mathbf{E}[Y]$. Solution: Since X and Y are positive, we have $\max(X, Y) \leq X + Y$. This gives us $\mathbf{E}[\max(X, Y)] = \mathbf{E}[X + Y]$ and by linearity of expectation, we then have $\mathbf{E}[X + Y] = \mathbf{E}[X] + \mathbf{E}[Y]$. Thus, $\mathbf{E}[\max(X, Y)] = \mathbf{E}[X] + \mathbf{E}[Y]$. \Box
- 3. Prove that $\operatorname{Var}[aX] = a^2 \operatorname{Var}[X]$, if a is a constant. **Proof:** Let Y = aX. Then $\operatorname{Var}[aX] = \operatorname{Var}[Y]$. Observe that

$$\begin{aligned} \mathbf{Var}[Y] &= \mathbf{E}[Y^2] - (\mathbf{E}[Y])^2 \ by \ definition \ of \ variance \\ &= \mathbf{E}[a^2 \cdot X^2] - (\mathbf{E}[a \cdot X])^2 \\ &= a^2 \cdot \mathbf{E}[X^2] - a^2 (\mathbf{E}[X])^2 \end{aligned}$$

By definition, $a^2 \cdot \mathbf{E}[X^2] - a^2 \cdot (\mathbf{E}[X])^2 = a^2 \cdot (\mathbf{E}[X^2] - (\mathbf{E}[X])^2) = a^2 \mathbf{Var}[X]$. Therefore, $\mathbf{Var}[aX] = a^2 \mathbf{Var}[X]$. \Box

4. Assume that X can take on only two values, viz. 0 and 1. Show that $\operatorname{Var}[X] = \mathbf{E}[X] \cdot \mathbf{E}[1 - X]$. Solution: Let $\operatorname{Pr}\{X = 0\} = 1 - p$ and let $\operatorname{Pr}\{X = 1\} = p$. Then, by definition, $\mathbf{E}[X] = 0 \cdot (1-p) + 1 \cdot p = p$ and $\mathbf{E}[X^2] = 0^2 \cdot (1-p) + 1^2 \cdot p = p$, thus $\mathbf{E}[X^2] = \mathbf{E}[X]$. Observe that, by definition, $\operatorname{Var}[X] = \mathbf{E}[X^2] - (\mathbf{E}[X])^2$. Since, $\mathbf{E}[X^2] = \mathbf{E}[X]$, we have

$$\begin{aligned} \mathbf{Var}[X] &= \mathbf{E}[X] - (\mathbf{E}[X])^2 \\ &= \mathbf{E}[X] \cdot (1 - \mathbf{E}[X]) \\ &= \mathbf{E}[X] \cdot \mathbf{E}[1 - X] \ by \ linearity \ of \ expectation \end{aligned}$$

Thus, $\mathbf{Var}[X] = \mathbf{E}[X] \cdot \mathbf{E}[1 - X]$. \Box

5. Let X be non-negative. Show that $\Pr\{X \ge t\} \le \frac{\mathbf{E}[X]}{t}$, for t > 0. (Markov's inequality.) **Proof:** Let $E[X] = \sum_{x \cdot p(x) > 0} x \cdot p(x)$. By splitting the summation, we obtain $\sum_{x \cdot p(x) > 0} x \cdot p(x) = \sum_{x \ge t} x \cdot p(x) + \sum_{x < t} x \cdot p(x)$. Using the technique of bounding terms, we then obtain

$$\begin{split} \sum_{x \geq t} x \cdot p(x) + \sum_{x < t} x \cdot p(x) \geq \sum_{x \geq t} x \cdot p(x) \\ \geq \sum_{x \geq t} t \cdot p(x) \end{split}$$

Since, $\sum_{x \ge t} t \cdot p(x) = t \cdot \sum_{x \ge t} p(x) = t \cdot \mathbf{Pr}\{x \ge t\}$, we have that $\mathbf{Pr}\{X \ge t\} \le \frac{\mathbf{E}[X]}{t}$. \Box

6. Let μ and σ denote the expectation and variance of X respectively. Prove that $\mathbf{Pr}\{|X - \mu| \ge t\sigma\} \le \frac{1}{t^2}$. (*Chebshev's inequality.*)

Proof: Observe that $(X - \mu)^2$ is a non-negative random variable, with expectation σ (as per the definition of variance), i.e., $\mathbf{E}[(X - \mu)^2] = \sigma^2$.

Therefore, we can use Markov's Inequality, to get $\mathbf{Pr}\{(X-\mu)^2 \ge \sigma^2 \cdot t^2\} \le \frac{\mathbf{E}[(X-\mu)^2]}{\sigma^2 \cdot t^2} = \frac{\sigma^2}{\sigma^2 \cdot t^2} = \frac{1}{t^2}$. Observe that since $(X-\mu)^2$ is non-negative, $(X-\mu)^2 \ge t^2 \sigma^2 \Leftrightarrow |X-\mu| \ge t \cdot \sigma$. Therefore, $\mathbf{Pr}\{|X-\mu| \ge t \cdot \sigma\} \le \frac{1}{t^2}$. \Box