

# HAMILTONIAN

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- Computationally very different from Eulerian paths
- Note that **HAMILTONIAN** is in **NP**
- We reduce 3-CNF SAT to **HAMILTONIAN**
  - ▶ For a fixed 3-CNF SAT formula, show that it can be transformed into a graph whose Hamiltonian path will give us the assignments for SAT.

## Reduction Sketch

$$\phi = \mathcal{C}_1 \wedge \mathcal{C}_2 \wedge \dots \wedge \mathcal{C}_m, \quad n \text{ variables}, \quad \mathcal{C}_i = (x_1^i \vee x_2^i \vee x_3^i)$$

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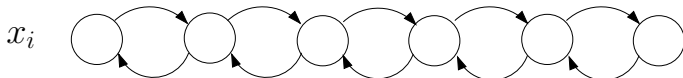
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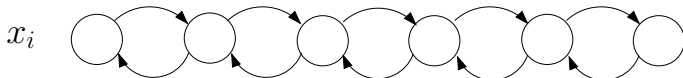
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- We first define the mapping of variables, and then the clauses.
- Each  $x_i$  will correspond to a path (chain) of  $6m$  vertices
- If we are at the first (or end) vertex, only one path to follow



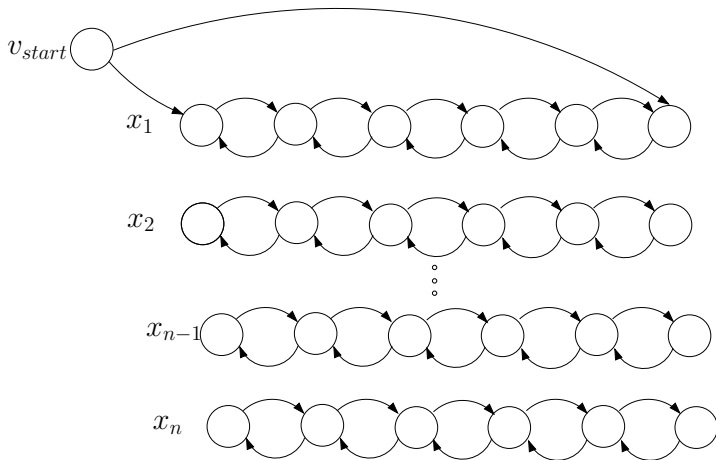
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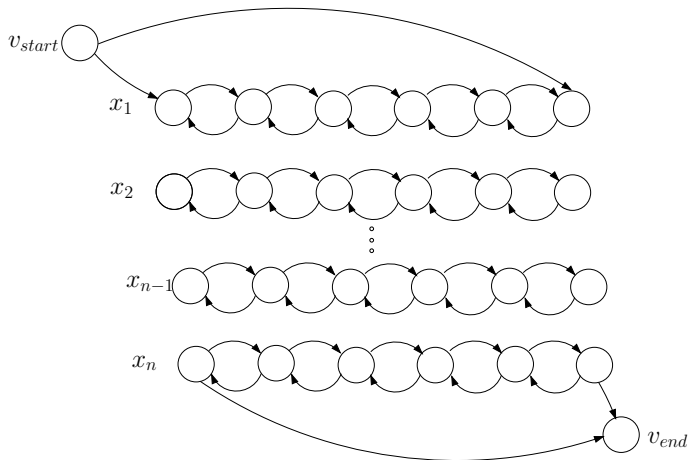
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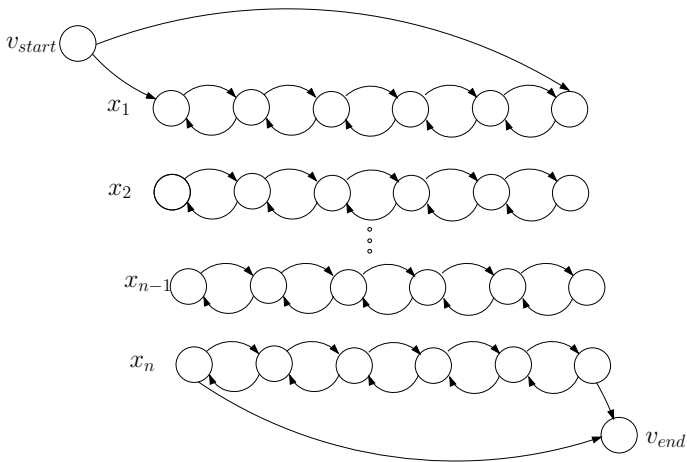
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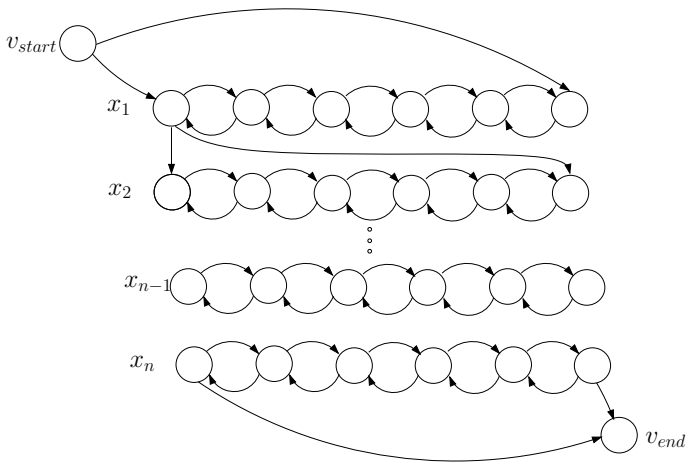
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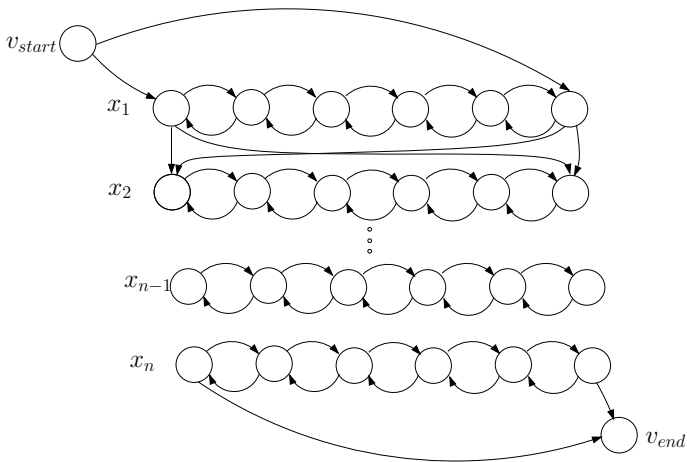
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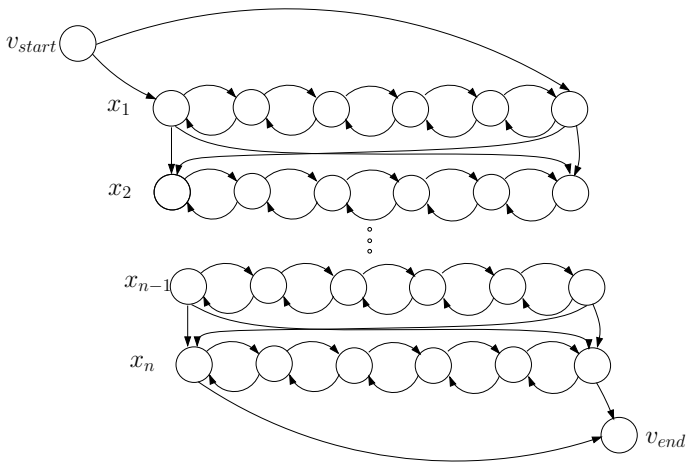
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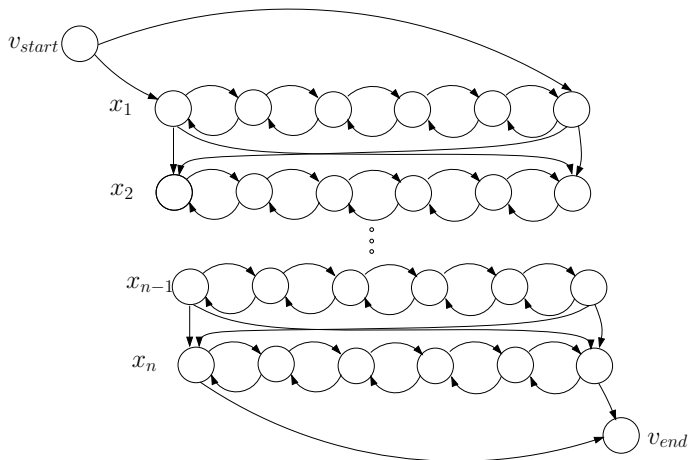
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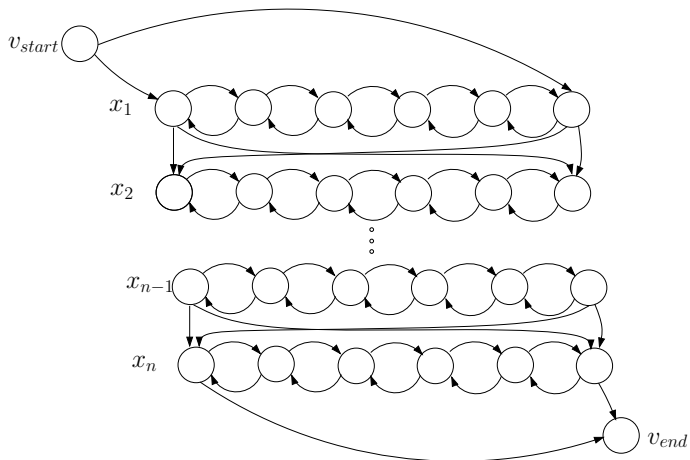


# Construction Properties



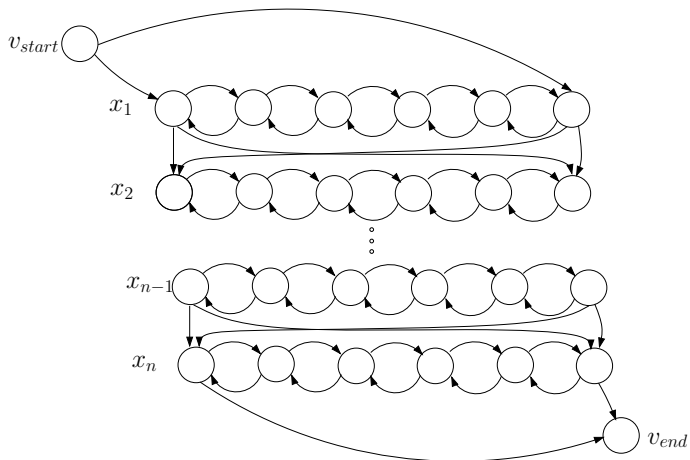
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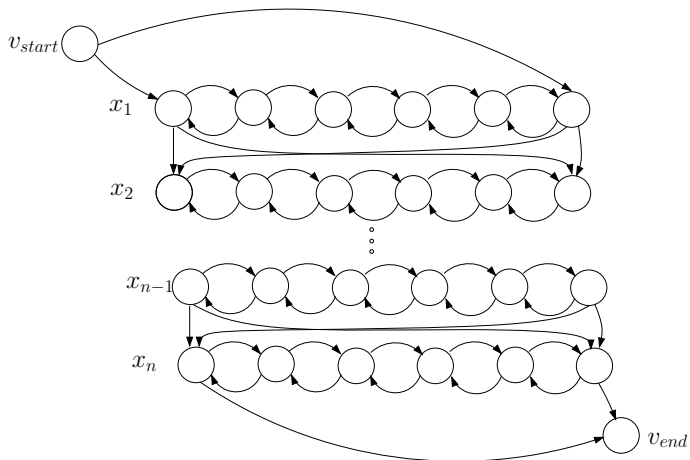
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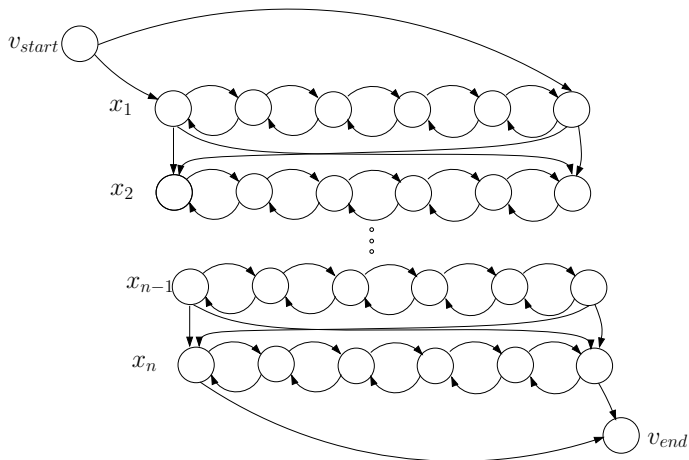
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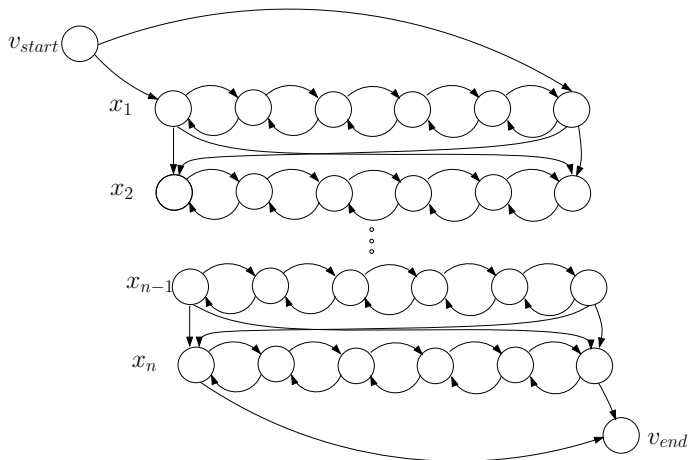
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- For each chain, only two ways of traversing it.
  - ▶ Left-to-right means  $x_i = 1$ , right-to-left means  $x_i = 0$



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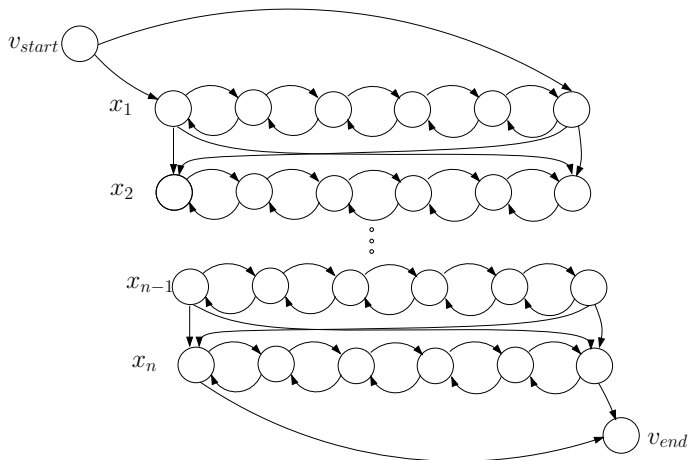
- Each assignment of variables corresponds to a unique Hamiltonian path.





# Construction Properties

- So far, no constraints – they will come from the clauses now.



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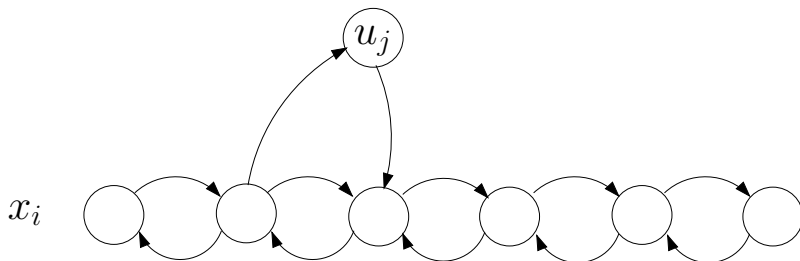
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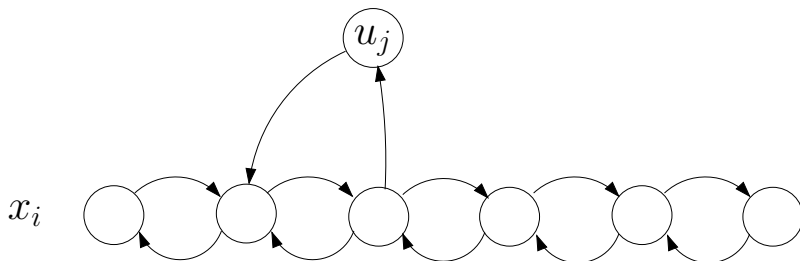
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## An Example

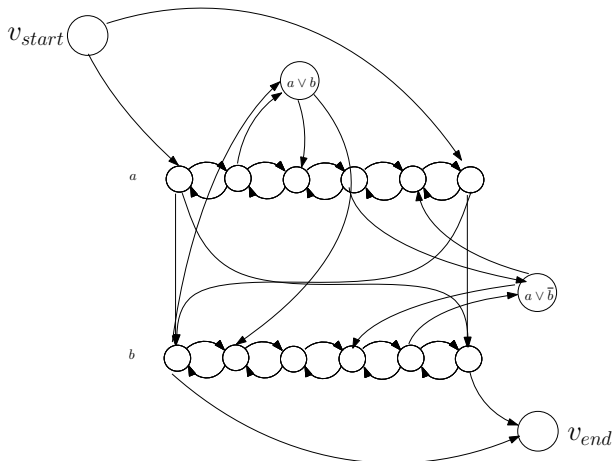
Construction of Hamiltonian path for

$$(a \vee b) \wedge (a \vee \bar{b})$$

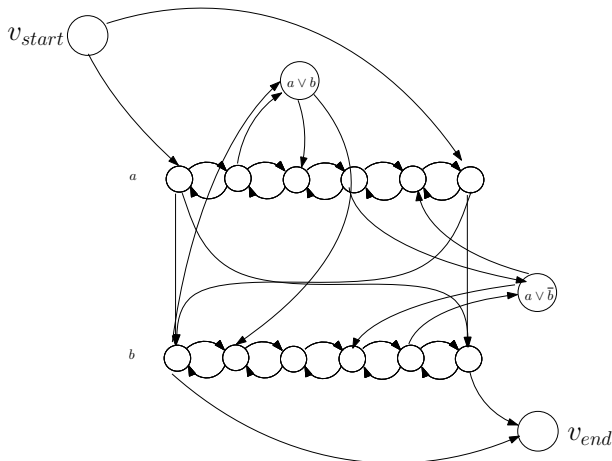
Here:

- $C_1 = (a \vee b)$
- $C_2 = (a \vee \bar{b})$

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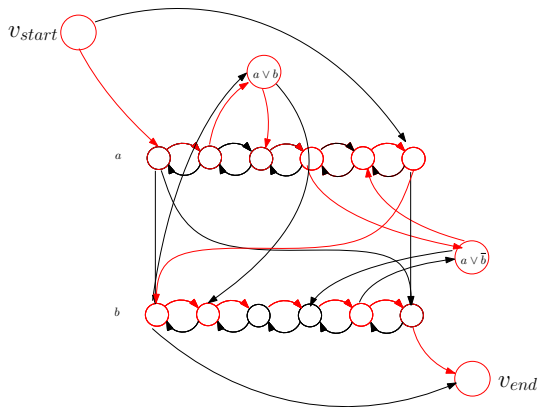
## Claim

Hamiltonian path exists **ONLY** if you go from left to right in ***a*** and chose any one of the two directions for ***b***.



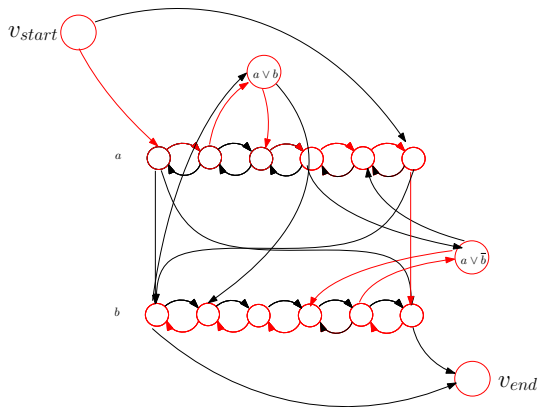
# Path Corresponding to Assignment 1

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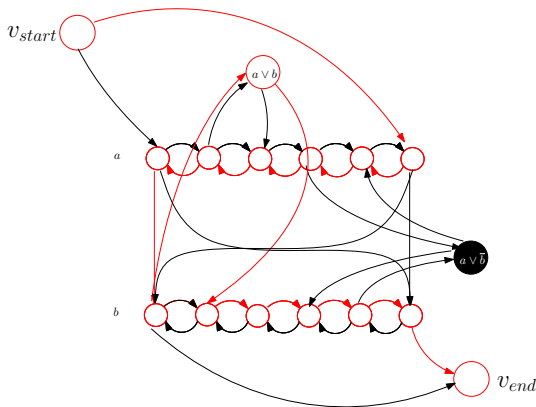
## Path Corresponding to Assignment 2

$a = 1; b = 0$



## Path Corresponding to Assignment 3

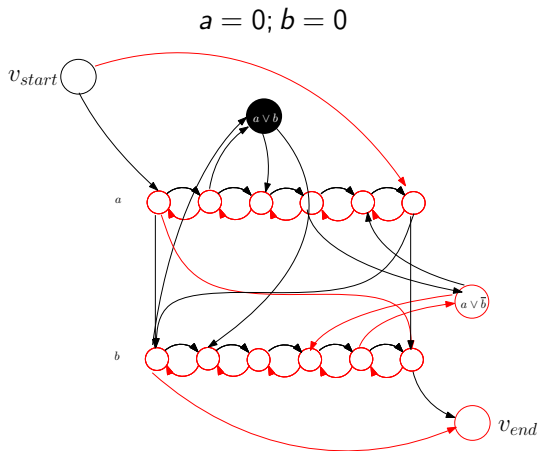
$$a = 0; b = 1$$



Error in finding HAMILTONIAN

No HAMILTONIAN PATH as  $(a \vee \bar{b})$  is not accessible

## Path Corresponding to Assignment 4



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- The above two *only* ways to visit  $u_j$  without getting stuck.