

Mathematical Preliminaries

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- Asymptotics and Inequalities
- Probability and Expectation





- Asymptotics and Inequalities
- Probability and Expectation

Abstract Algebra





- Asymptotics and Inequalities
- Probability and Expectation
- Abstract Algebra
- 5 Upper and Lower Bounds





- Asymptotics and Inequalities
- Probability and Expectation
- Abstract Algebra
- 5 Upper and Lower Bounds
- Problem paradigms





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🕜 Roadmap

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Alphabets and strings

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Alphabets and strings

Basics

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Alphabets and strings



Alphabet -

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Alphabet - A finite, non-empty collection of symbols,

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- Alphabet A finite, non-empty collection of symbols, e.g., $\Sigma = \{a, b, ...\}$.
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Languages and Problems Asymptotics and Inequalities Probability and Expectation

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- Powers of an alphabet Σ^k is the set of all strings of length k, each of whose symbols is in Σ.

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$$\Sigma^* = \cup_{i=0}^{\infty} \Sigma^i$$

Languages

Asymptotics and Inequalities Probability and Expectation Abstract Algebra Upper and Lower Bounds Problem paradigms Roadmap

Languages

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Any set $L \subseteq \Sigma^*$ is called a language over Σ .

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Give a language *L* and a string $w \in \Sigma^*$, decide whether $w \in L$.

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Primality can be thought of as membership in the language L_p , where L_p is the set of all binary strings whose value is a prime.

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Note

Languages and problems are two sides of the same coin.

Order of magnitude of functions

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(i) Which function grows faster: $100x^2$ or $\frac{1}{10^6}x^3$?

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- (i) Which function grows faster: $100x^2$ or $\frac{1}{10^6}x^3$?
- (ii) Which function grows faster: $x^2 10$ or x + 10?

Order of Magnitude (contd.)

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Definition

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Let *f* and *g* be functions mapping non-negative reals to non-negative reals.

Order of Magnitude (contd.)

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Let *f* and *g* be functions mapping non-negative reals to non-negative reals.

Then f = O(g), if there exist constants c and n_0 such that for all $n \ge n_0$, $f(x) \le c \cdot g(x)$.

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Let *f* and *g* be functions mapping non-negative reals to non-negative reals.

Then f = o(g), if there exist constants c and n_0 such that for all $n \ge n_0$, $f(x) < c \cdot g(x)$.

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Then $f = \Theta(g)$, if f = O(g) and g = O(f).

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Examples

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(i) Let
$$f(x) = 2x^2 - 2$$
 and $g(x) = \frac{1}{100}x^2 - 100$.

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Test to determine order

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The limit test

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Let f and g denote two functions mapping non-negative reals to non-negative reals.

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Let f and g denote two functions mapping non-negative reals to non-negative reals.

Let
$$I = \lim_{x \to \infty} \frac{f(x)}{g(x)}$$
. Then

(i) If I is a positive constant,

Test to determine order

The limit test

Let f and g denote two functions mapping non-negative reals to non-negative reals.

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Let f and g denote two functions mapping non-negative reals to non-negative reals.

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(ii) If $I = 0$.

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Note

If
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 and if $\lim_{x\to\infty} g(x) = \infty$, then,

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The above rule is called L'Hospital's rule.

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(i) Show that
$$x = o(x^2)$$

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- (i) Show that $x = o(x^2)$.
- (ii) Show that $x = o(x \log x)$.

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- (i) Show that $x = o(x^2)$.
- (ii) Show that $x = o(x \log x)$.
- (iii) Show that $\log x = o(x)$.

Some inequalities

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Useful relationships

 $\bigcirc \log n! = \Theta(n \cdot \log n).$

Some inequalities

- $log n! = \Theta(n \cdot \log n).$
- **2** $n! = o(n^n), \omega(2^n).$

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$$n! = \sqrt{2 \cdot \pi \cdot n} \cdot (\frac{n}{e})^n \cdot (1 + \Theta(\frac{1}{n})).$$

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$$e^x = \sum_{i=0}^{\infty} \frac{x^i}{i!}$$

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$$e^{x} = \sum_{i=0}^{\infty} \frac{x^{i}}{i!}$$
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• $\ln(1+x) = \sum_{i=1}^{\infty} (-1)^{i+1} \cdot \frac{x^{i}}{i}$, when $|x| < 1$.

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$$\log n! = \Theta(n \cdot \log n).$$

• $n! = o(n^n), \omega(2^n).$

$$1 = \sqrt{2 \cdot \pi \cdot n} \cdot (\frac{n}{e})^n \cdot (1 + \Theta(\frac{1}{n})).$$

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• $\ln(1+x) = \sum_{i=1}^{\infty} (-1)^{i+1} \cdot \frac{x^{i}}{i!}$, when $|x| < 1$.

$$Iim_{n\to\infty} \frac{n^b}{a^n} = 0, \text{ for all } a > 1.$$

Some inequalities

Useful relationships

log n! = Θ(n ⋅ log n).
n! = o(nⁿ), ω(2ⁿ).
n! = √2 ⋅ π ⋅ n ⋅ (ⁿ/_e)ⁿ ⋅ (1 + Θ(¹/_n)).
e^x = ∑[∞]_{i=0} xⁱ/_{i!}.
ln(1 + x) = ∑[∞]_{i=1}(-1)ⁱ⁺¹ ⋅ xⁱ/_i, when |x| < 1.
lim_{n→∞} n^b/_{aⁿ} = 0, for all a > 1.
a = b^{log_b a}.

Some inequalities

Useful relationships

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Quick tricks through calculus

Quick tricks through calculus

Bounding sums through integration

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• Let *f* be a monotonically increasing function.
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Exercise

Find upper and lowee bounds on $\sum_{i=1}^{n} \frac{1}{i}$.

Sample Space and Events

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Subramani Computational Complexity

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- (i) Suppose that the experiment consists of tossing a coin. Then, $S = \{H, T\}$.
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Any subset of the sample space S is called an event.

Combining Events

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Subramani Computational Complexity

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Roadmap

Combining events (contd.)

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Definition

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Note

Never forget that events are sets.

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Note

Never forget that events are sets. This is particularly important when using logic to reason about them.

Languages and Problems Asymptotics and Inequalities Probability and Expectation Abstract Algebra

Abstract Algebra Upper and Lower Bounds Problem paradigms Roadmap

Defining Probabilities on Events

Problem paradigms Roadmap

Defining Probabilities on Events

Assigning probabilities

Roadmap

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Let *S* denote a sample space. We assume that the number P(E) is assigned to each event *E* in *S*, such that:

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Two Identities

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Conditional Probability

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Motivation

Subramani Computational Complexity

Conditional Probability

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In the previously discussed coin tossing example, let E denote the event that both coins turn up heads and F denote the event that the first coin turns up heads.

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Subramani Computational Complexity

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Consider the experiment of tossing two fair dice. Let F denote the event that the first die turns up 4. Let E_1 denote the event that the sum of the faces of the two dice is 6.

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Bayes' Formula

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Derivation

Subramani Computational Complexity

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Bayes' Formula

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Bayes' Formula

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= $P(E | F)P(F) + P(E | F^{c})P(F^{c})$

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$$P(E) = P(EF) + P(EF^c)$$

$$= P(E \mid F)P(F) + P(E \mid F^{c})P(F^{c})$$

$$= P(E \mid F)P(F) + P(E \mid F^c)(1 - P(F))$$

Bayes' Formula

Derivation

Let *E* and *F* denote two arbitrary events on a sample space *S*. Clearly, $E = EF \cup EF^c$, where the events *EF* and *EF*^{*c*} are mutually exclusive. Now, observe that,

$$P(E) = P(EF) + P(EF^{c}) = P(E | F)P(F) + P(E | F^{c})P(F^{c}) = P(E | F)P(F) + P(E | F^{c})(1 - P(F))$$

Thus, the probability of an event E

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Thus, the probability of an event *E* is the weighted average of the conditional probability of *E*, given that event *F* has occurred and the conditional probability of *E*, given that event *F* has not occurred,

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Thus, the probability of an event *E* is the weighted average of the conditional probability of *E*, given that event *F* has occurred and the conditional probability of *E*, given that event *F* has not occurred, each conditional probability being given as much weight as the probability of the event that it is conditioned on, has of occurring.

Random Variables

Random Variables

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Subramani Computational Complexity

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In case of certain random experiments, we are not so much interested in the actual outcome, but in some function of the outcome, e.g., in the experiment of tossing two dice, we could be interested in knowing whether or not the the sum of the upturned faces is 7. We may not care whether the actual outcome is $(1, 6), (6, 1), or \ldots$

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Upper and Lower Bounds Problem paradigms Roadmap

The Bernoulli Random Variable

Problem paradigms Roadmap

The Bernoulli Random Variable

Main idea

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If we let the random variable X assume the value 1, if the experiment was a success and 0, if the experiment was a failure, then X is said to be a Bernoulli random variable.

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 $p(1) = P\{X = 1\} = p$ $p(0) = P\{X = 0\} = 1 - p.$ Languages and Problems Asymptotics and Inequalities Probability and Expectation Abstract Algebra

Upper and Lower Bounds Problem paradigms Roadmap

The Binomial Random Variable

Problem paradigms Roadmap

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Motivation

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The probability mass function of *X* is given by:

$$p(i) = P\{X = i\} = C(n, i) \cdot p^{i} \cdot (1 - p)^{n-i}, i = 0, 1, 2, \dots n$$

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The Geometric Random Variable

Problem paradigms Roadmap

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Roadmap

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Problem paradigms Roadmap

Features of a random variable

Roadmap

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Distribution Function

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Definition (Distribution Function)

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For a random variable X, the distribution function $F(\cdot)$ is defined for any real number b, $-\infty < b < \infty$, by

 $F(b) = P(X \leq b).$

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Parameters of the important Random Variables

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Parameter table

Variable type	Expectation	Variance
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Binomial	n · p	$n \cdot p \cdot (1-p)$
Geometric	$\frac{1}{p}$	$\frac{1-p}{p^2}$

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Find the parameters of the Poisson, Normal, Uniform and exponential random variables.

Expectation of the function of a random variable

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Theorem

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Upper and Lower Bounds Problem paradigms Roadmap

Independent Random Variables

Roadmap

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 $Var(X_1 + X_2) = Var(X_1) + Var(X_2) + 2 \cdot Cov(X_1, X_2).$

Concentration Inequalities

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Tail bounds

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Inequality	Known parameters	Tail bound
Markov	$X \ge 0, E[X]$	$P(X \ge a \cdot E[X]) \le \frac{1}{a}, \ a > 0$
Chebyshev	E[X], Var(X)	$P(X - E[X] \ge a \cdot E[X]) \le rac{Var(X)}{(a \cdot E[X])^2}, a > 0.$
Chernoff	X is binomial, $E[X]$	$P((X - E[X]) \ge \delta) \le e^{-\frac{-2\cdot\delta^2}{n}}, \delta > 0.$

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Chebyshev	E[X], Var(X)	$P(X - E[X] \ge a \cdot E[X]) \le \frac{\operatorname{Var}(X)}{(a \cdot E[X])^2}, a > 0.$
Chernoff	X is binomial, $E[X]$	$P((X - E[X]) \ge \delta) \le e^{-\frac{-2\cdot\delta^2}{n}}, \delta > 0.$

Exercise

Find the tail bounds for the coin tossing problem using all three techniques.

Problem paradigms Roadmap

Groups

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Subramani Computational Complexity

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• Associativity: $(a \cdot b) \cdot c = a \cdot (b \cdot c)$, for all $a, b, c \in G$.

2 Identity: There exists an element $e \in G$ such that $a \cdot e = e \cdot a = a$ for all $a \in G$.

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Example

 \mathcal{Z} with + as the operator and 0 as the identity element.

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Upper and Lower Bounds Problem paradigms Roadmap

Rings

Problem paradigms Roadmap

Rings

Definition

Subramani Computational Complexity

Rings

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- **()** Left and right distributivity: For all $a, b, c \in R$, we have, $a \cdot (b + c) = a \cdot b + a \cdot c$ and $(b + c) \cdot a = b \cdot a + c \cdot a$.

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- **(**) Multiplicative associativity: For all $a, b, c \in R$, we have, $(a \cdot b) \cdot c = a \cdot (b \cdot c)$.

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- **6** Multiplicative associativity: For all $a, b, c \in R$, we have, $(a \cdot b) \cdot c = a \cdot (b \cdot c)$.

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Example

 $\ensuremath{\mathcal{Z}}$ under the usual addition and multiplication.

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Upper and Lower Bounds Problem paradigms Roadmap

Fields

Roadmap

Fields

Definition

Subramani Computational Complexity

Fields

Definition

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Definition

A field consists of a set F and two binary operations "+" (addition) and " \cdot " (multiplication), defined on R, for which the following conditions are satisfied:

• $(F, +, \cdot)$ is a ring.

Fields

Definition

- (F, +, \cdot) is a ring.
- 2 Multiplicative commutative:

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A field consists of a set F and two binary operations "+" (addition) and " \cdot " (multiplication), defined on R, for which the following conditions are satisfied:

• (F, +, \cdot) is a ring.

2 Multiplicative commutative: For any $a, b \in F$, $a \cdot b = b \cdot a$.

Fields

Definition

- (F, +, \cdot) is a ring.
- 2 Multiplicative commutative: For any $a, b \in F$, $a \cdot b = b \cdot a$.
- O Multiplicative identity:

Fields

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- (F, +, \cdot) is a ring.
- 2 Multiplicative commutative: For any $a, b \in F$, $a \cdot b = b \cdot a$.
- **3** Multiplicative identity: There exists $1 \in F$ such that $a \cdot 1 = 1 \cdot a = a$ for all $a \in F$.

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Example

The set \Re with traditional addition and multiplication.

Bound Computation

Bound Computation

Upper and Lower bounds

Subramani Computational Complexity

Bound Computation

Upper and Lower bounds

Maximum/Minimum in an array.

Bound Computation

Upper and Lower bounds

- Maximum/Minimum in an array.
- 2 Maximum and minimum in an array.

Bound Computation

Upper and Lower bounds

- Maximum/Minimum in an array.
- 2 Maximum and minimum in an array.
- Array sorting.

Bound Computation

Upper and Lower bounds

- Maximum/Minimum in an array.
- 2 Maximum and minimum in an array.
- Array sorting.
- Matrix multiplication.

Problem types

Problem types

Three types of problems

Subramani Computational Complexity

Problem types

Three types of problems

• The Königsberg bridge problem.

Subramani Computational Complexity

Problem types

Three types of problems

- The Königsberg bridge problem.
- 2 The Hamilton Circuit problem.

Problem types

Three types of problems

- The Königsberg bridge problem.
- 2 The Hamilton Circuit problem.
- O Playing Chess.

The path forward

The path forward

Roadmap

Subramani Computational Complexity

The path forward

Roadmap

Some problems permit insights.
The path forward

- Some problems permit insights.
- One problems have short proofs.

The path forward

- Some problems permit insights.
- One problems have short proofs.
- One power of programming languages.

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- Memory issues.

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- Some problems permit insights.
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- Coin tossing and adversary.

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