Robust Optimization: Applications

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Outline



1 Overview

- Theory of Robust Optimization
- Applications of Robust Optimization

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- Applications of Robust Optimization

2 A

Applications

- Shortest Path Problem
- Facility Location Problem
- Portfolio Selection Problem

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Reasoning

- The data of real-world optimization problems often are not known exactly at the time the problem is being solved.
- In real-world applications of optimization, even a small uncertainty in the data can make the nominal optimal solution to the problem completely meaningless from a practical viewpoint.
- Consequently, in optmization, there exist a real need of a methodology capable of detecting cases when data uncertainty can heavily effect the quality of the nominal solution.

Note

SH. Amini Optimization Methods in Finance

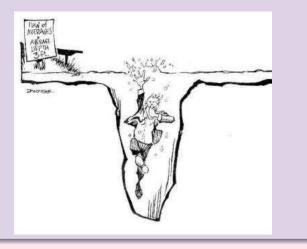
Note

• Optimization based on nominal values often lead to SEVERE infeasibilities.

Applications of Robust Optimization

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SH. Amini	Optimization Methods in Finance
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 $3.030 \cdot 1000 - 2.999 \cdot 1000 = 31 > 1$

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 - Often difficult to specify reliably the distribution of uncertain data.

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Robust Optimization:

- Does not assume stochastic nature of the uncertain data.
- Remains computationally tractable.
- Sometimes is too pessimistic.

LP Programming with uncertainty

LP Programming with uncertainty

• Standard LP form:

 $\begin{aligned} \min z &= \mathbf{c} \cdot \mathbf{x} \\ \mathbf{A} \cdot \mathbf{x} &\leq \mathbf{b} \end{aligned}$

LP Programming with uncertainty

• Standard LP form:

$$\min z = \mathbf{c} \cdot \mathbf{x}$$

L

• Uncertain LP form:

$$\min z = \widetilde{\mathbf{c}} \cdot \mathbf{x}$$
$$\widetilde{\mathbf{A}} \cdot \mathbf{x} \leq \widetilde{\mathbf{b}}$$

• Certain Linear Constrain:

 $\textbf{A}\cdot\textbf{x} \quad \leq \quad \textbf{b}$

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$$\widetilde{\mathbf{A}} \cdot \mathbf{x} \leq \mathbf{b}$$

• Chance Constrain:

$$\mathbf{P}(\widetilde{\mathbf{A}} \cdot \mathbf{x} > \mathbf{b}) \leq \mathcal{E}$$

Affine Uncertainy

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$$\widetilde{\mathbf{A}} = \mathbf{A}(\widetilde{\mathbf{z}}) = \mathbf{A}^{\mathbf{0}} + \sum_{j=1}^{N} \mathbf{A}_{j} \cdot \widetilde{\mathbf{z}}_{j}$$

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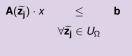
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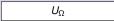
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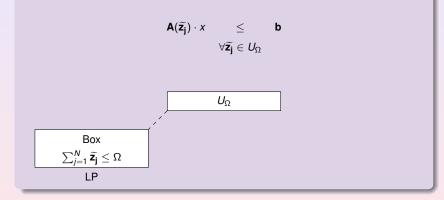
$$(\mathbf{A^0} + \sum_{j=1}^N \mathbf{A_j} \cdot \widetilde{\mathbf{z}_j}) \cdot x \le \mathbf{b}$$

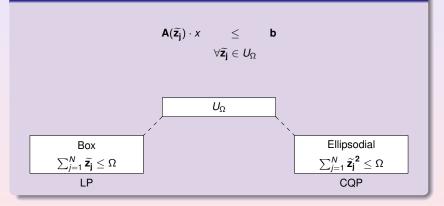
Theory of Robust Optimization Applications of Robust Optimization

$$egin{array}{lll} \mathsf{A}(\widetilde{\mathsf{z}_{\mathsf{j}}})\cdot x &\leq \mathsf{b} \ & orall \widetilde{\mathsf{z}_{\mathsf{i}}}\in U_{\Omega} \end{array}$$









• Measure of the level of robustness.

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- Used to adjust conservativeness of the robust optimization solution.

 $\Omega \in (0,\sqrt{N})$

Applications of Robust Optimization

Budget of Uncertainty (Ω)

- Measure of the level of robustness.
- Used to adjust conservativeness of the robust optimization solution.

$$\Omega \in (0,\sqrt{N})$$

• $\Omega = 0$ means no uncertainty and $\Omega = \sqrt{N}$ is worse case budget.

$$\mathcal{E} = exp(-\frac{\Omega^2}{2})$$

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$$\mathcal{E} = exp(-\frac{0^2}{2}) = 1.000$$

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 - Production scheduling: uncertainty of grade and price;
 - Mineral processing circuit design: uncertainty of grade, feed rate and price.

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• Circuit design: minimizing delay in digital circuits when the underlying gate delays are not known exactly.

- Mining engineering:
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• Civil Engineering:

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- Facility locations: uncertainty of demand.

• Electrical Engineering:

- Circuit design: minimizing delay in digital circuits when the underlying gate delays are not known exactly.
- Finance:
 - Robust portfolio optimization: random returns;
 - Robust risk management: uncertainty of Value at Risk.

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Overview Applications Shortest Path Problem Facility Location Problem Portfolio Selection Problem

Example

Given a weighted diagraph, find the shortest path from 1 to 3 if:

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• Link travel costs (c_j) are subject to bounded uncertainty;

Example

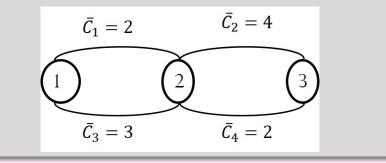
Given a weighted diagraph, find the shortest path from 1 to 3 if:

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- There are two sources of uncertainty (\tilde{z}_i) .

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Overview Applications Shortest Path Problem Facility Location Problem Portfolio Selection Problem

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Example

$\widetilde{\textbf{c_1}}=2+1.5\widetilde{z_1}+1.5\widetilde{z_2}$

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$$\widetilde{\mathbf{c_1}} = 2 + 1.5\widetilde{z_1} + 1.5\widetilde{z_2}$$
$$\widetilde{\mathbf{c_2}} = 4 + 0.5\widetilde{z_1} + 0.5\widetilde{z_2}$$

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$$\widetilde{\mathbf{c}_3} = 3 + 0.5\widetilde{z_1} + 0.5\widetilde{z_2}$$

Example

$$\begin{split} \widetilde{\mathbf{c_1}} &= 2 + 1.5 \widetilde{z_1} + 1.5 \widetilde{z_2} \\ \widetilde{\mathbf{c_2}} &= 4 + 0.5 \widetilde{z_1} + 0.5 \widetilde{z_2} \\ \widetilde{\mathbf{c_3}} &= 3 + 0.5 \widetilde{z_1} + 0.5 \widetilde{z_2} \\ \widetilde{\mathbf{c_4}} &= 2 + 1.5 \widetilde{z_1} + 1.5 \widetilde{z_2} \end{split}$$

Example

$$\widetilde{\mathbf{c}_1} = 2 + 1.5\widetilde{z_1} + 1.5\widetilde{z_2}$$
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where:

 $\widetilde{z_1}^2 + \widetilde{z_2}^2 \leq \Omega$ $\Omega \in (0, 1, 2)$

Ω=0	\tilde{y}_{k1}	\tilde{y}_{k1}	ROBUST SP	CLASSICAL SP
1-2	0	0	6	6
1-4	0	0	4	4
3-2	0	0	7	7
3-4	0	0	5	5

Ω=1	\tilde{y}_{k1}	\tilde{y}_{k1}	ROBUST SP	CLASSICAL SP
1-2	$\sqrt{2}/2$	$\sqrt{2}/2$	8.83	6
1-4	$\sqrt{2}/2$	$\sqrt{2}/2$	8.24	4
3-2	$\sqrt{2}/2$	$\sqrt{2}/2$	9.82	7
3-4	$\sqrt{2}/2$	$\sqrt{2}/2$	7.82	5

Ω=2	\tilde{y}_{k1}	\tilde{y}_{k1}	ROBUST SP	CLASSICAL SP
1-2	1	1	10.00	6
1-4	1	1	10.00	4
3-2	1	1	11.00	7
3-4	1	1	9.00	5

Shortest Path Problem

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2 Applications

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Application of Facility Location Models

• New airport, hospital, and school.

Application of Facility Location Models

- New airport, hospital, and school.
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- New airport, hospital, and school.
- Addition of a new workstation.
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Application of Facility Location Models

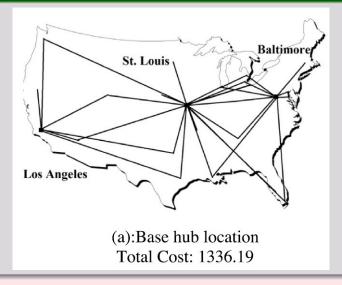
- New airport, hospital, and school.
- Addition of a new workstation.
- Warehouse location.
- Bathroom location in a facility etc..

Example

• The goal of the facility location problem for airline industry is to find an optimal location of hubs.

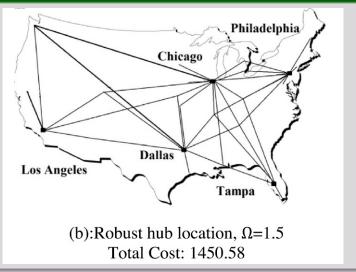
- The goal of the facility location problem for airline industry is to find an optimal location of hubs.
- Where demand is uncertain and its distribution is not fully specified.

Overview overview Shortest Path Problem Facility Location Problem Portfolio Selection Problem



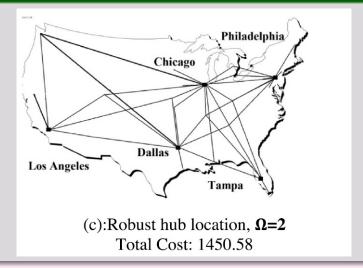
Overview overview Shortest Path Problem Facility Location Problem Portfolio Selection Problem

Example



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Shortest Path Problem

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Example

We consider the following simple portfolio optimization example:

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Maximize $\mu_1 \cdot x_1 + \mu_2 \cdot x_2 + \mu_3 \cdot x_3$

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subject to

 $\begin{array}{rcl} TE(x_1, x_2, x_3) &\leq & 0.10 \\ x_1 + x_2 + x_3 &= & 1 \\ x_1, x_2, x_3 &\geq & 0 \end{array}$

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 x_1 and x_2 are first and second assets.

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The benchmark is the portfolio that invests funds half-and-half in the two assets.

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 x_3 represents proportion of the funds that are not invested.

The benchmark is the portfolio that invests funds half-and-half in the two assets.

TE(x) represents the tracking error of the portfolio with respect to the half-and-half benchmark.

Example

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$$TE = \sqrt{ \begin{bmatrix} x_1 - 0.5 \\ x_2 - .05 \\ x_3 \end{bmatrix}^T \begin{bmatrix} 0.1764 & 0.09702 & 0 \\ 0.09702 & 0.1089 & 0 \\ 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 - 0.5 \\ x_2 - .05 \\ x_3 \end{bmatrix}}$$

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Example

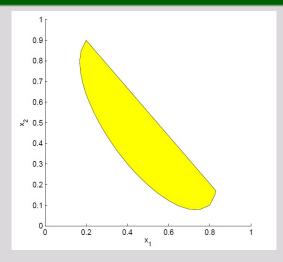


Figure: The feasible set of the portfolio selection problem

Example

We now build a relative robustness model for this portfolio problem:

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• Scenario 1: (μ_1, μ_2, μ_3) : (6, 4, 0)

Example

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- Scenario 1: (μ₁, μ₂, μ₃) : (6, 4, 0)
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So, the objective values will be:

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So, the objective values will be:

• Scenario 1: 5.662

Example

We now build a relative robustness model for this portfolio problem:

- Scenario 1: (μ₁, μ₂, μ₃) : (6, 4, 0)
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- Scenario 3: (μ_1, μ_2, μ_3) : (4, 6, 0)

So, the objective values will be:

- Scenario 1: 5.662
- Scenario 2: 5.662

Example

We now build a relative robustness model for this portfolio problem:

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So, the objective values will be:

- Scenario 1: 5.662
- Scenario 2: 5.662
- Scenario 3: 5.000

Example

Relative robust formulation:

Example

Relative robust formulation:

 $\min_{x,t} t$

Example

$$\min_{\substack{x,t} \\ 5.662 - (6x_1 + 4x_2) \leq t$$

Example

$$\min_{x,t} t$$
5.662 - (6x₁ + 4x₂) $\leq t$
5.662 - (4x₁ + 6x₂) $< t$

Example

$$\begin{array}{rcl} \min_{x,t} & t \\ 5.662 - (6x_1 + 4x_2) & \leq & t \\ 5.662 - (4x_1 + 6x_2) & \leq & t \\ 5.000 - (5x_1 + 5x_2) & \leq & t \end{array}$$

Example

$$\begin{array}{rcl} \min_{x,t} & t \\ 5.662 - (6x_1 + 4x_2) & \leq & t \\ 5.662 - (4x_1 + 6x_2) & \leq & t \\ 5.000 - (5x_1 + 5x_2) & \leq & t \\ TE(x_1, x_2, x_3) & \leq & 0.10 \end{array}$$

Example

$$\begin{array}{rcl} & \min_{x,t} & t \\ 5.662 - (6x_1 + 4x_2) & \leq & t \\ 5.662 - (4x_1 + 6x_2) & \leq & t \\ 5.000 - (5x_1 + 5x_2) & \leq & t \\ TE(x_1, x_2, x_3) & \leq & 0.10 \\ x_1 + x_2 + x_3 & = & 1 \end{array}$$

Example

$$\begin{array}{rcl} \min_{x,t} & t \\ 5.662 - (6x_1 + 4x_2) &\leq t \\ 5.662 - (4x_1 + 6x_2) &\leq t \\ 5.000 - (5x_1 + 5x_2) &\leq t \\ TE(x_1, x_2, x_3) &\leq 0.10 \\ x_1 + x_2 + x_3 &= 1 \\ x_1, x_2, x_3 &\geq 0 \end{array}$$

Example

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choosing a maximum tolerable regret level of 0.75 we get the following feasibility problem:

Find x

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choosing a maximum tolerable regret level of 0.75 we get the following feasibility problem:

Find x $5.662 - (6x_1 + 4x_2) \leq 0.75$

Example

F	ind	x	
$5.662 - (6x_1 + 4)$	x ₂)	\leq	0.75
$5.662 - (4x_1 + 6)$	<i>x</i> ₂)	\leq	0.75

Example

Find	x	
$5.662 - (6x_1 + 4x_2)$	\leq	0.75
$5.662 - (4x_1 + 6x_2)$	\leq	0.75
$5.000 - (5x_1 + 5x_2)$	\leq	0.75

Example

Find	X	
$5.662 - (6x_1 + 4x_2)$	\leq	0.75
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$5.000 - (5x_1 + 5x_2)$	\leq	0.75
$TE(x_1, x_2, x_3)$	\leq	0.10

Example

Find	X	
$5.662 - (6x_1 + 4x_2)$	\leq	0.75
$5.662 - (4x_1 + 6x_2)$	\leq	0.75
$5.000 - (5x_1 + 5x_2)$	\leq	0.75
$TE(x_1, x_2, x_3)$	\leq	0.10
$x_1 + x_2 + x_3$	=	1

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$x_1 + x_2 + x_3$	=	1
x_1, x_2, x_3	\geq	0

Overview Applications Shortest Path Problem Facility Location Problem Portfolio Selection Problem

Example

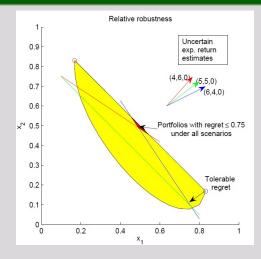


Figure: Set of solutions with regret less than 0.75

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